

Crossroads 1.77

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Abstract

Crossroads is a load balance and fail over utility for TCP based services. It is a daemon program running in user space, and features extensive configurability, polling of back ends using 'wakeup calls', detailed status reporting, 'hooks' for special actions when backend calls fail, and much more. Crossroads is service-independent: it is usable for HTTP/HTTPS, SSH, SMTP, DNS, etc. In the case of HTTP balancing, Crossroads can modify HTTP headers, e.g. to provide 'session stickiness' for back end processes that need sessions, but aren't session-aware of other back ends. Crossroads is configured via a file and accessible via a web interface.

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1 Introduction

Crossroads is a daemon that basically accepts TCP connections at preconfigured ports, and given a list of 'back ends' distributes each incoming connection to one of the back ends, so that a client request is served. Additionally, crossroads maintains an internal administration of the back end connectivity: if a back end isn't usable, then the client request is handled using another back end. Crossroads will then periodically check whether a previously not usable back end has come to life yet. Also, crossroads can select back ends by estimating the load, so that balancing is achieved.

Using this approach, crossroads serves as load balancer and fail over utility. Crossroads will very likely not be as reliable as hardware based balancers, since it always will require a server to run on. This server, in turn, may become a new Single Point of Failure (SPOS). However, in situations where cost efficiency is an issue, crossroads may be a good choice. Furthermore, crossroads can be deployed in situations where a hardware based balancing already exists and augmenting service reliability is needed. Or, crossroads may be run off a diskless system, which again improves reliability of the underlying hardware.

This document describes how to use crossroads, how to configure it in order to increase the reliability of your systems, and how to compile the program from its sources.

1.1 Obtaining Crossroads

As quick reference, here are some important URL's for Crossroads:

- <http://crossroads.e-tunity.com> is the site that serves Crossroads. You can browse this at leisure for documentation, sources, and so on.
- <http://freshmeat.net/projects/crossr> is the Freshmeat announcement page.
- [svn://svn.e-tunity.com/crossroads](http://svn.e-tunity.com/crossroads) is the SVN repository; anonymous reading (fetching) is allowed. In order to commit changes, mail the author or maintainer for credentials.

1.2 Reporting bugs

Crossroads was thoroughly tested and proven to work. However, on particular Unices with particular needs and network particularities, bugs may occur.

In such cases you can contact the maintainer to ask for assistance. Visit the site <http://crossroads.e-tunity.com> for contact information. In questions or bug reports, always include the following:

- A description of the bug: when does it occur, what are the circumstances, what is the expected behavior, what is the observed behavior.
- The Crossroads configuration (normally `/etc/crossroads.conf`)
- The version of Crossroads and all compile-time settings, obtained via `crossroads -C`.

1.3 Copyright and Disclaimer

Crossroads is distributed as-is, without assumptions of fitness or usability. You are free to use crossroads to your liking. It's free, and as with everything that's free: there's also no warranty. Crossroads is distributed under the GNU General Public Licence, version 3. See <http://crossroads.e-tunity.com> for more information.

You are allowed to make modifications to the source code of crossroads, and you are allowed to (re)distribute crossroads, as long as you include this text, all sources, and if applicable: all your modifications, with each distribution.

While you are allowed to make any and all changes to the sources, I would appreciate hearing about them. If the changes concern new functionality or bugfixes, then I'll include

them in a next release, stating full credits. If you want to seriously contribute (to which you are heartily encouraged), then mail me and I'll get you access to the Crossroads SVN repository, so that you can update and commit as you like.

1.4 Terminology

Throughout this document, the following terms are used: ¹

A client is a process that initiates a network connection to get contact with some service.

A service or server process or listener is a central application that accepts network connections from clients and services them.

Back ends are locations where crossroads looks in order to service its clients. Crossroads sits 'in between' and does its tricks. Therefore, as far as the back ends are concerned, crossroads behaves like a client. As far as the true client is concerned, crossroads behaves like the service. The communication is however transparent: neither client nor back end are aware of the middle position of crossroads.

A connection is a network conversation between client and service, where data are transferred to and fro. As far as crossroads is concerned, success means that a connection can be established without errors on the network level. Crossroads isn't aware of service peculiarities. E.g., when a webserver answers `HTTP/1.0 500 Server Error` then crossroads will see this as a succesful connection, though the user behind a browser may think otherwise.

Back end selection algorithms are methods by which crossroads determines which back end it will talk to next. Crossroads has a number of built-in algorithms, which may be configured per service.

Back end states are the statuses of each back end that is known to crossroads. A back end may be available, (temporarily) unavailable or truly down. When a back end is temporarily unavailable, then crossroads will periodically check whether the back end has come to life yet (that is, if configured so).

A spike is a sudden increase in activity, leading to extra load on a given service. When crossroads is in effect and when the spike occurs in one connection, then obviously the spike will also appear at one of the back ends. However, crossroads will see the spike and will make sure that a subsequent request goes to an other back end. In contrast, when several connections arrive simultaneously and cause a spike, then crossroads will be able to distribute the connections over several back ends, thereby 'flattening out' the increase.

Load balancing means that incoming client requests are distributed over more than just one back end (which wouldn't be the case if you wouldn't be running crossroads). Enabling load balancing is nothing more than duplicating services over more than one back end, and having something (in this case: crossroads) distribute the requests, so that per back end the load doesn't get too high.

An HTTP session is a series of separate network connections that originate from one browser. E.g., to fill the display with text and images, the browser hits a website several times. An HTTP session may even span several screens. E.g., a website registration dialog may involve 3 screens that when called from the same browser, form a logical group of some sort.

Headers or header lines are specific parts of an HTTP message. Crossroads has directives to add or modify headers that are part of the request that a browser sends to server, or those that are part of the server.

¹Many more meanings of the terms will exist - yes, I am aware of that. I'm using the terms here in a very strict sense.

Session stickiness means that when a browser starts an HTTP dialog, the balancer makes sure that it 'sticks' to the same back end (i.e., subsequent requests from the browser are forced to go to the same back end, instead of being balanced to other ones).

Back end usage is measured by crossroads in order to be able to determine back end selection. Crossroads stores information about the number of active connections, the transferred bytes and about the connection duration. These numbers can be used to estimate which back end is the least used – and therefore, presumably, the best candidate for a new request.

Fail over is almost always used when load balancing is in effect. The distributor of client requests (crossroads of course) can also monitor back ends, so that incase a back end is 'down', it is no longer accessed.

Service downtime normally occurs when a service is switched off. Downtime is obviously avoided when fail over is in effect: a back end can be taken out of service in a controlled manner, without any client noticing it.

1.5 Porting Issues

This section lists some caveats when converting Crossroads configurations to new versions. Given the changes of the syntax of the configuration file of Crossroads, existing configuration files may need to be made suitable for new versions.

1.5.1 Porting issues for pre-1.63 installations

The parser of Crossroads 1.63 has become even stricter. HTTP-related commands, such as `addclientheader`, may only occur when the service type is `http`.

Previously such directives were allowed, albeit they would be ineffective.

1.5.2 Porting issues for pre-1.50 installations

As of version 1.50, the command `crossroads tell` has been changed. To set a state, the command must be used as follows:

```
crossroads tell service backend state newstate
```

The keyword `state` was added. This is because `crossroads tell` can also set a new server address; in that case the keyword `server` is required. If you have automatic health checkers or the like that change the back end states, please modify the commands to suit this.

1.5.3 Porting issues for pre-1.43 installations

As of version 1.43, the shared memory key calculations are based on a different algorithm. This key is e.g. necessary when starting and stopping Crossroads; in both actions, the key must be computed in the same way.

Therefore, when upgrading Crossroads, make sure that you stop a running Crossroads daemon using the same binary that started it. After this, install a new binary, and start the daemon using the new binary.

(Incidentally, this is always a good idea, but especially so when the 'old' binary is pre-1.43 and the 'new' one is post-1.43.)

Furthermore, shell-style comment is no longer supported as of 1.43. The reason is that 1.43 introduces the hash mark as a preprocessor token (in `#include`, `#define`). Therefore, if your configuration files use shell-style comment, please convert this to **C** or **C++** style.

1.5.4 Porting issues for pre-1.21 installations

As of version 1.21, the event-hook directives `onsuccess` and `onfailure` no longer exists.

- Please replace `onsuccess` by `onstart`;
- Please replace `onfailure` by `onfail`;
- Note that there is a new hook `onend`.

The commands that are run via `onstart`, `onend` or `onfail` are subject to format expansion; e.g., `%1w` is expanded to the weight of the first back end, etc.. See section 4 for details.

1.5.5 Porting issues for pre-0.26 installations

As of version 0.26 the syntax of the configuration file has changed. In particular:

- The keyword `maxconnections` is now used instead of `maxclients`;
- The keyword `connectiontimeout` is now used instead of `sessiontimeout`.

Therefore when converting configuration files to the new syntax, the above keywords must be changed. (The reason for these changes is that 0.26 introduces *sticky HTTP sessions* that span multiple TCP connections, and the term *session* is used strictly in that sense – and no longer for a TCP connection.)

1.5.6 Porting issues for pre-1.08 installations

As of version 1.08, the following directives no longer are supported:

- `insertstickycookie` was replaced by the more generic directive `addclientheader`. E.g., instead of
`insertstickycookie "XRID=100; Path=/"`;
the syntax is now
`addclientheader "Set-Cookie: XRID=100; Path=/"`;
- `insertrealip` was replaced by the more generic directive `setserverheader`. E.g., instead of
`insertrealip on`;
the syntax is now
`setserverheader "XR-Real-IP: %r"`;
This incidentally also makes it possible to change the header name (here: `XR-Real-IP`).

2 Installation for the impatient

For the impatient, here's the very-quick-but-very-superficial recipe for getting Crossroads up and running:

- If you don't have SVN or don't want to use it:
 - Obtain the Crossroads source archive at <http://crossroads.e-tunity.com>.
 - Change-dir to a 'sources' directory on your system and unpack the archive.
 - Change-dir into the create directory `crossroads/`.
- If you have SVN and want to go for the newest snapshot:
 - Get the latest sources and snapshots using SVN from `svn://svn.e-tunity.com/crossroads`.
 - You'll find the newest alpha version under `crossroads/trunk` and the stable versions under `crossroads/tags`, e.g. `crossroads/tags/release-1.00`.

- Choose which you want to use: the latest stable release, or the bleeding edge alpha? In the former case, change-dir to `crossroads/tags/release-X.YY`, where `X.YY` is a release ID. In the latter case, change-dir to `crossroads/trunk`.
- Type `make install`. This installs the Crossroads binary into `/usr/local/bin/`. If the compilation doesn't work on your system, check `etc/Makefile.def` for hints.
- Create a file `/etc/crossroads.conf`. In it state something like:

```
service www {
    port 80;
    revivinginterval 15;
    backend one {
        server 10.1.1.100:80;
    }
    backend two {
        server 10.1.1.101:80;
    }
}
```

That's off course assuming that you want to balance HTTP on port 80 to two back ends at 10.1.1.100 and 10.1.1.101.

- Type `crossroads start`.
- Surf to the machine where Crossroads is running. You will see the pages served by the back ends 10.1.1.100 or 10.1.1.101.
- To monitor the status of Crossroads from the command line, type `crossroads status`.
- If you want to monitor Crossroads via a web front end, start the manager: type `crossroads-mgr start 1000`. Next point your browser to the machine where Crossroads is running, port 1000.

3 Using Crossroads

3.1 The Balancer: crossroads

The Crossroads balancer is started from the commandline, and highly depends on `/etc/crossroads.conf` (the default configuration file). It supports a number of flags (e.g., to overrule the location of the configuration file). The actual usage information is always obtained by typing `crossroads` without any arguments. Crossroads then displays the allowed arguments.

The installation (see section 7) installs also a second program called `crossroads-daemon`. This program is not meant to be started from the command line; it is administered through the front end `crossroads`.

3.1.1 General Commandline Syntax

This section shows the most basic usage. As said above, start `crossroads` without arguments to view the full listing of options.

- `crossroads start` and `crossroads stop` are typical actions that are run from system startup scripts. The meaning is self-explanatory.
- `crossroads restart` is a combination of the former two. Beware that a restart may cause discontinuity in service; it is just a shorthand for typing the 'stop' and 'start' actions after one another.
- `crossroad status` reports on each running service. Per service, the state of each back end is reported.

- `crossroads tell service backend state newstate` is a command line way of telling crossroads that a given back end, of a given service, is in a given state. Normally crossroads maintains state information itself, but by using `crossroads tell`, a back end can be e.g. taken 'off line' for servicing.
- `crossroads tel service backend server hostname:port` redefines the back end address of a running back end.
- `crossroads configtest` tells you whether the configuration is syntactically correct.
- `crossroads services` reports on the configured services. In contrast to `crossroads status`, this option only shows what's configured – not what's up and running. Therefore, `crossroads services` doesn't report on back end states.

3.1.2 Status Reporting from the Command Line

The command `crossroads status` shows a verbose human-readable report of how Crossroads is doing. When many services are configured, this can be a somewhat lengthy output. If you're interested in the overview of only one service, you can use `crossroads status servicename`, in which case the report will only be shown for the stated service.

Similarly, if you're interested only in the status of a given back end of a given service, use `crossroads service servicename backendname`.

The flag `-t` causes the status overview to be abbreviated and displayed in a parseable format. This flag can be used in automated scripts that check how Crossroads is doing; e.g., in health checking scripts. When `-t` is used, the format of the status reporter is as follows:

- Service health is reported on one line per service.
- The first string on the line is the service name.
- After this, a series follows of `backendname=state`, where the back end name is the configured name of the back end, and where state is e.g. `available`, `down`, `unavailable`, `waking`. The statuses may occur in caps. The series is repeated for all back ends of the given service.

The flag `-x` causes the overview to be presented as an XML document. This can be handy if you want to further automate the control over Crossroads, or if you want to show the status via the web. (See section 5.15 for more information.)

The status reporter shows the number of running connections per back end, and the total of connections of the entire service. Due to technical reasons the total count can be one off. (There is a technical reason for this: the total counter gets updated at a different rate than separate back end counts. This may be fixed in a future release.)

3.1.3 Logging-related options

Two 'flags' of Crossroads are specifically logging-related. This section elaborates on these flags.

First, there's flag `-a`. When present, the start and end of activity is logged using statements like

YYYY-MM-DD HH/MM/SS starting http from 61.45.32.189 to 10.1.1.1

Similarly, there are 'ending' statements. Using this flag and scanning your logs for these statements may be helpful in quickly determining your system load.

Second, there's flag `-l`. This flag selects the 'facility' of logging and defaults to `LOG_DAEMON`. You can supply a number between 0 and 7 to flag `-l` to select `LOG_LOCAL0` to `LOG_LOCAL7`. This would separate the Crossroads-related logging from other streams. Here's a very short guide; please read your Unix manpages of `syslogd` for more information.

- First edit `/etc/syslog.conf` and add a line:

```
local7.*    /var/log/crossroads.log
```

That instructs `syslogd` to send `LOG_LOCAL7` requests to the logfile `/var/log/crossroads.log`.

- Next, restart `syslogd`. On most Unices that's done by issuing `killall -1 syslogd`. (As a side-note, I tried this once on an Bull/AIX system, and the box just shut down. The `killall` command killed every process...)
- Now start `crossroads` with the flag `-l7`.
- Finally, monitor `/var/log/crossroads.log` for Crossroads' messages.

3.1.4 Reloading Configurations

Crossroads doesn't support the reloading of a configuration while running (such as other programs, e.g. Apache do). There are various technical reasons for this.

However, external lists of allowed or denied IP addresses can be reloaded by sending a signal `-1` (`SIGHUP`) to Crossroads. See section 4.3 for the details.

3.2 The Web Frontend: `crossroads-mgr`

A web front end for Crossroads is `crossroads-mgr`. Type `crossroads-mgr` without arguments to see the usage.

Normally the web interface is started like:

```
crossroads-mgr start 1000
```

where 1000 is a free TCP port. Note that `crossroads-mgr` must be started on the system where the balancer `crossroads` itself is running.

There are lots of other options that can be used with `crossroads-mgr`, see section 5.14.

4 The configuration

The configuration that `crossroads` uses is normally stored in the file `/etc/crossroads.conf`. This location can be overruled using the command line flag `-c`.

This section explains the syntax of the configuration file, and what all settings do.

4.1 General language elements

This section describes the general elements of the `crossroads` configuration language.

4.1.1 Empty lines, indentation and comments

Empty lines are of course allowed in the configuration. Also, indentation is irrelevant as far as processing is concerned, but is loosely inspired on a C-like style. So the following two examples are identical as far as Crossroads is concerned (though the first one will be typically more readable):

```
/* Example one */
options {
    tcpbuffersize 10240;
    logactivity on;
}

/* Example two */
options {tcpbuffersize 10240
;logactivity on;}
```

Crossroads recognizes the following comment formats:

- C-style, between `/*` and `*/`,
- C++-style, starting with `//` and ending with the end of the text line.

Simply choose your favorite editor and use the comment that 'looks best'.²

4.1.2 Preprocessor directives

Similar to C or C++, the Crossroads grammar knows `#include` and `#define`. Both directives must start at the first column of the line (ie., the `#` sign must occur at the leftmost line position).

- `#include "filename"` includes the stated file name at the place of the statement.
- `#define SYMBOL DEFINITION` defines `SYMBOL` as placeholder for `DEFINITION`.

For example, one may use the configuration:

```
#define SERVICEPORT 80
service web {
    port SERVICEPORT;
    .
    . /* More statements follow here */
    .
}
```

The port statement is then read as `port 80`.

The statement `#define` can also be very nicely used when trying out Crossroads configurations. Crossroads has a statement `verbosity true` that causes debugging information to be logged. Once a configuration has proven to work, you'll most likely want `verbosity false` so that overhead due to logging is avoided. This can be easily implemented using `#define`:

```
/* Set DEBUG to true or false;
 * true   is for testing purposes,
 * false  is for production */
#define DEBUG true

service web {
    verbosity DEBUG;
    .
    . /* More statements follow here */
    .
}
```

4.1.3 Keywords, numbers, identifiers, generic strings

In a configuration file, statements are identified by *keywords*, which are reserved words. The list of the keywords of the Crossroads grammar is: `service port verbosity maxconnections type any http backend server bindto connectiontimeout on off dispatchmode roundrobin random byduration bysize byconnections byorder byclientip over decay revivinginterval checkinterval retries shmkey weight onstart onfail backlog throughput log traffic log http timing log stickycookie addclientheader setclientheader appendclientheader addserverheader setserverheader appendserverheader allowfrom denyfrom allowfile denyfile externalhandler useraccount onend headerinspection deep shallow state available unavailable down options logactivity tcpbuffer size dnscache ttl logfacility shmpermissions sloppyportbind leaveprocesstitle`

²I favor C or C++ comment. My favorite editor *emacs* can be put in `cmode` and nicely highlight what's comment and what's not. And as a bonus it will auto-indent the configuration!

Many keywords require an *identifier* as the argument. E.g., a service has a unique name, which must start with a letter or underscore, followed by zero or more letters, underscores, or digits. E.g., in the statement `service myservice`, the keyword is `service` and the identifier is `myservice`.

Other keywords require a numeric argument. Crossroads knows only non-negative integer numbers, as in `port 8000`. Here, `port` is the keyword and `8000` is the number. Octal numbers are specified by prefixing a zero; e.g., `0644` is an octal number.

Yet other keywords require 'generic strings', such as hostname specifications or system commands. Such generic strings contain any characters (including white space) up to the terminating statement character `;`. If a string must contain a semicolon, then it must be enclosed in single or double quotes:

- `This is a string;` is a string that starts at `T` and ends with `g`
- `"This is a string";` is the same, the double quotes are not necessary
- `"This is ; a string";` has double quotes to protect the inner `;`

Finally, an argument can be a 'boolean' value. Crossroads knows `true`, `false`, `yes`, `no`, `on`, `off`. The keywords `true`, `yes` and `on` all mean the same and can be used interchangeably; as can the keywords `false`, `no` and `off`.

4.2 Daemon options

Crossroads supports an optional block that defines run options for the daemon process. These options can also be specified on the command line using flags.

The syntax of the options block is:

```
options {  
    .  
    . option statements (see below)  
    .  
}
```

In the curly-brace block the following option statements can occur:

- `logactivity on;` or `off` Turns activity logging on or off. This is also controlled using the flag `-a`. The default is not to log network activity.
- `tcpbuffersize number ;` This defines the size of network buffers. The default is 5120 bytes. This is also controlled using the flag `-B`.
- `dnscahettl number ;` This controls the time-to-live of cached DNS entries, in seconds. Value zero means no caching and is the default. This is also controlled using the flag `-d`.
- `logfacility number ;` This controls the logging facility. Values 0 to 7 select `LOG_LOCAL0` to `LOG_LOCAL7`. Any other value selects `LOG_DAEMON`, which is also the default. This is also controlled using the flag `-l`.
- `shmpermissions number ;` This defines the permissions (for user, group and other) of the shared memory block. The number is most often specified as an octal value, e.g. `0644`, which is also the default. This is also controlled using the flag `-m`.
- `sloppyportbind on;` or `off` Turns on 'sloppy' port binding of the listener. When `on`, Crossroads will not treat port-busy conditions as fatal, but will wait and retry. The default is `off`. This is also controlled using flag `-s`.
- `leaveprocesstitle on;` or `off` When `on`, the process title of Crossroads is not modified, so that `ps` will show `crossroads-daemon`. The default is `off`: Crossroads will modify its process name into something like `crossroads - Service web: listening` or `crossroads - Service web: serving`.

4.3 Service definitions

Service definitions are blocks in the configuration file that state what is for each service. A service definition starts with `service`, followed by a unique identifier, and by statements in `{` and `}`. For example:

```
// Definition of service 'www':
service www {
    ...
    ... // statements that define the
    ... // service named 'www'
    ...
}
```

The configuration file can contain many service blocks, as long as the identifying names differ. The following list shows possible statements. Each statement must end with a semicolon, except for the `backend` statement, which has its own block (more on this later).

4.3.1 port - Specifying the listen port

Description: The `port` statement defines to which TCP port a service 'listens'. E.g. `port 8000` says that this service will accept connections on port 8000.

Multiple services in one configuration cannot use the same port number, unless they both bind to specific (and different) IP addresses. See also the `bindto` statement.

Syntax: `port number`

Default: There is no default. This is a required setting.

4.3.2 type - Defining the service type

Description: The `type` statement defines how crossroads handles the stated service. There are currently two types: `any` and `http`. The type `any` means that crossroads doesn't interpret the contents of a TCP stream, but only distributes streams over back ends. The type `http` means that crossroads has to analyze what's in the messages, does magical HTTP header tricks, and so on – all to ensure that multiple connections are treated as one session, or that the back end is notified of the client's IP address.

Unless you really need such special features, use the type `any` (the default), even for HTTP protocols.

Syntax: `type specifier`, where *specifier* is `any` or `http`

Default: `any`

4.3.3 headerinspection - are all HTTP headers inspected

Description: The `headerinspection` directive defines whether Crossroads must inspect all HTTP headers that are seen on one TCP connection, or only the first ones. There are two possible values for this directive: `deep` and `shallow`. In `deep` mode, all information that is seen on the TCP link is monitored and parsed, and all HTTP header blocks are analyzed and subject to directives such as `addclientheader`. In `shallow` mode, only the first header block that the server sends, and the first header block that forms the server's answer, are analyzed.

Syntax: `headerinspection specifier`, where *specifier* is `deep` or `shallow`

Default: `deep`

4.3.4 bindto - Binding to a specific IP address

Description: The `bindto` statement is used in situations where crossroads should only listen to the stated port at a given IP address. E.g., `bindto 127.0.0.1` causes crossroads to 'bind' the service only to the local IP address. Network connections from other hosts won't be serviced. By default, crossroads binds a service to all presently active IP addresses at the invoking host.

Syntax: `bindto address`, where *address* is a numeric IP address, such as 127.0.0.1, or the keyword `any`.

Default: `any`

4.3.5 verbosity - Controlling debug output

Description: Verbosity statements come in two forms: `verbosity on` or `verbosity off`. When 'on', log messages to `/var/log/messages` are generated that show what's going on.³ The keyword `verbose` is an alias for `verbosity`.

Syntax: `verbosity setting` or `verbose setting`, where *setting* is `true`, `yes` or `on` to turn verbosity on; or `false`, `no`, `off` to turn it off.

Default: `off`

4.3.6 dispatchmode - How are back ends selected

Description: The dispatch mode controls how crossroads selects a back end from a list of active back ends. The below text shows the bare syntax. See section 5.2 for a textual explanation. The settings can be:

- `dispatchmode roundrobin`: Simply the 'next in line' is chosen. E.g, when 3 back ends are active, then the usage series is 1, 2, 3, 1, 2, 3, and so on. Roundrobin dispatching is the default method, when no `dispatchmode` statement occurs.
- `dispatchmode random`: Random selection. Probably only for stress testing, though when used with weights (see below) it is a good distributor of new connections too.
- `dispatchmode bysize [over connections]`: The next back end is the one that has transferred the least number of bytes. This selection mechanism assumes that the more bytes, the heavier the load. The modifier `over connections` is optional. (The square brackets shown above are not part of the statement but indicate optionality.) When given, the load is computed as an average of the last stated number of connections. When this modifier is absent, then the load is computed over all connections since startup.
- `dispatchmode byduration [over connections]`: The next back end is the one that served connections for the shortest time. This mechanism assumes that the longer the connection, the heavier the load.
- `dispatchmode byconnections`: The next back end is the one with the least active connections. This mechanism assumes that each connection to a back end represents load. It is usable for e.g. database connections.
- `dispatchmode byorder`: The first back end is selected every time, unless it's unavailable. In that case the second is taken, and so on.

³Actually, the messages go to `syslog(3)`, using facility `LOG_DAEMON` and priority `LOG_INFO`. In most (Linux) cases this will mean: output to `/var/log/messages`. On Mac OSX the messages go to `/var/log/system.log`.

- `dispatchmode byclientip`: The client's IP address is 'hashed' into a number, which is used to pick a back end. The same client IP address will therefore always be dispatched to the same back end. When the back end of choice is down, `dispatchmode byconnections` is used.
- `dispatchmode externalhandler` *program arguments*: This is a special mode, where an external program is delegated the responsibility to say which back end should be used next. In this case, Crossroads will call the external program, and this will of course be slower than one of the 'built-in' dispatch modes. However, this is the ultimate escape when custom-made dispatch modes are needed. The dispatch mode that uses an `externalhandler` is discussed separately in section 5.11.

The selection algorithm is only used when clients are serviced that aren't part of a sticky HTTP session. This is the case during:

- all client requests of a service type `any`;
- new sessions of a service type `http`.

When type `http` is in effect and a session is underway, then the previously used back end is always selected – regardless of dispatching mode.

Your 'right' dispatch mode will depend on the type of service. Given the fact that crossroads doesn't know (and doesn't care) how to estimate load from a network traffic stream, you have to choose an appropriate dispatch mode to optimize load balancing. In most cases, `roundrobin` or `byconnections` will do the job just fine.

Syntax: `dispatchmode mode` (see above for the modes), optionally followed by `over number`, or when the `mode` is `externalhandler`, followed by *program*.

Default: `roundrobin`

4.3.7 `revivinginterval` - Back end wakeup calls

Description: A reviving interval definition is used when Crossroads determines that a back end is temporarily unavailable. This will happen when:

- The back end cannot be reached (network connection fails);
- The network connection to the back end suddenly dies.

Once a reviving interval is set, Crossroads will periodically check the unavailable back end(s) to see whether they have woken up.

An example of the definition is `revivinginterval 10`. When this reviving interval is given, crossroads will check each 10 seconds whether unavailable back ends have woken up yet.

Syntax: • `revivinginterval number`;
 • `revivinginterval number externalhandler program arguments`;

The first form connects to a back end server. If the connection succeeds, then the back end is considered available. The second form activates an external program (see section 5.3 for a description). The back end is marked available if the program's exit status is zero.

Default: 0 (no wakeup calls)

4.3.8 `checkinterval` - Periodic back end checks

Description: When a check interval is stated, Crossroads will periodically probe back ends to determine whether available back ends are still there, and to see whether unavailable back ends have woken up yet.

An example is `checkinterval 10`. When this is stated, Crossroads will probe all back ends each 10 seconds.

Syntax: • `checkinterval number`;
 • `checkinterval number externalhandler program arguments`;

The first form checks by connecting to the back end server. If the connection succeeds, then the back end is considered available; otherwise the back end is considered unavailable.

The second form activates an external program (see section 5.3 for a description). The back end is considered available if the program's exit status is zero; otherwise it is considered unavailable.

Default: 0 (no periodic checks)

4.3.9 maxconnections - Limiting concurrent clients at service level

Description: The maximum number of connections is specified using `maxconnections`. There is one argument; the number of concurrent established connections that may be active within one service.

'Throttling' the number of connections is a way of preventing Denial of Service (DOS) attacks. Without a limit, numerous network connections may spawn so many server instances, that the service ultimately breaks down and becomes unavailable.

Note that `maxconnections` is also allowed in a backend description block, in which case it limits the number of TCP connections to that particular back end.

Syntax: `maxconnections number`, where the number specifies the maximum of concurrent connections to the service.

Default: 0, meaning that all connections will be accepted.

4.3.10 backlog - The TCP Back Log size

Description: The TCP back log size is a number that controls how many 'waiting' network connections may be queued, before a client simply cannot connect. The syntax is e.g. `backlog 5` to cause crossroads to have 5 waiting connections for 1 active connection. The backlog queue shouldn't be too high, or clients will experience timeouts before they can actually connect. The queue shouldn't be too small either, because clients would be simply rejected. Your mileage may vary.

Syntax: `backlog number`

Default: 0, which takes the operating system's default value for socket back log size.

4.3.11 shmkey - Shared Memory Access

Description: Different Crossroads invocations must 'know' of each others activity. E.g, `crossroad status` must be able to get to the actual state information of all running services. This is internally implemented through shared memory, which is reserved using a key.

Normally crossroads will supply a shared memory key, based on the service name. In situations where this conflicts with existing keys (of other programs, having their own keys), you may supply a chosen value.

The actual key value doesn't matter much, as long as it's unique and as long as each invocation of crossroads uses it.

Syntax: `shmkey number`

Default: 0, which means that crossroads will 'guess' its own key, based on the service name.

4.3.12 allow* and deny* - Allowing or denying connections

Description: Crossroads can allow or deny connections based on the IP address of a client. There are four directives that are relevant: `allowfrom`, `allowfile`, `denyfrom` and `denyfile`.

When using `allowfrom` and `denyfrom` then the IP addresses to allow or deny connections are stated in `/etc/crossroads.conf`. When using `allowfile` and `denyfile` the allow or deny connections are stated in a separate file.

When `allow*` directives are used, then all connections are denied unless they match the stated allowed IP's. When `deny*` directives are used, then all connections are allowed unless they match the stated disallowed IP's. When denying and allowing is both used, then the Crossroads checks the deny list first.

The statements `allowfrom` and `denyfrom` are followed by a list of filter specifications. The statements `allowfile` and `denyfile` are followed by a filename; Crossroads will read filter specifications from those external files. In both cases, Crossroads obtains filter specifications and places them in its lists of allowed or denied IP addresses. The difference between specifying filters in `/etc/crossroads.conf` or in external files, is that Crossroads will reload the external files when it receives signal 1 (SIGHUP), as in `killall -1 crossroads`.

The filter specifications must obey the following syntax: they are series of space-separated strings, consisting of up to four numbers ranging from 0 to 255 and separated by a decimal sign. Optionally a slash follows, with a bitmask which is also a decimal number. For example, `127.0.0/24 10/8 192.168.2.1` is a setting that consists of three specifiers. This is probably best explained by a few examples:

- `allowfrom 10/8`; will allow connections from `10.*.*.*` (a full Class A network). The mask `/8` means that the first 8 bits of the number (ie., only the 10) are significant. On the last 3 positions of the IP address, all numbers are allowed. Given this directive, client connections from e.g. `10.1.1.1` and `10.2.3.4` will be allowed.
- `allowfrom 10.3/16`; will allow all IP addresses that start with `10.3`. The first 16 bits (i.e., the first 2 numbers) are significant, the rest doesn't matter.
- `allowfrom 10.3.1/16`; is the same as above. The third byte of the IP address is superfluous because the netmask specifies that only the first 16 bits (2 numbers) are taken into account.
- `allowfrom 10.3.1.15`; allows traffic from only the specified IP address. There is no bitmask; all four numbers are relevant.
- `allowfrom 10.3.1.15 10.2/16`; allows traffic from one IP address `10.3.1.15` or from a complete Class B network `10.2.*.*`
- `allowfile /tmp/myfile.txt`; in combination with a file `/tmp/myfile.txt`, with the contents `10.3.1.15 10.2/16`, is the same as above.

When using `ttt(allowfrom)` and `denyfrom`, separate specifiers can be stated in one statement (separated by whitespace), or the whole statement can be repeated. E.g., the following two alternatives have the same effect:

```
/* Alternative 1: */
allowfrom 10/8 192.168.1/24;

/* Alternative 2: */
allowfrom 10/8;
allowfrom 192.168.1.24;
```

Syntax:

- `allowfrom filter-specification(s)`
- `denyfrom filter-specification(s)`

- `allowfile filename`
- `denyfile filename`

Default: In absence of these statements, all client IP's are accepted.

4.3.13 `useraccount` - Limiting the effective ID of external processes

Description: Using the directive `useraccount`, the effective user and group ID can be restricted. This comes into effect when Crossroads runs external commands, such as:

- Hooks for `onstart`, `onend` or `onfail`;
- External dispatchers, when `dispatchmode externalhandler` is in effect.

Once a user name for external commands is specified, Crossroads assumes the associated user ID and group ID before running those commands.

Syntax: `useraccount username`

Default: None; when unspecified, external commands are run with the ID that was in effect when Crossroads was started.

4.4 Backend definitions

Inside the service definitions as are described in the previous section, *backend definitions* must also occur. Backend definitions are started by the keyword `backend`, followed by an identifier (the back end name), and statements inside `{` and `}`:

```
service myservice {
    ...
    ... // statements that define the
    ... // service named 'myservice'
    ...

    backend mybackend {
        ...
        ... // statements that define the
        ... // backend named 'mybackend'
        ...
    }
}
```

Each service definition must have at least one backend definition. There may be more (and probably will, if you want balancing and fail over) as long as the backend names differ. The statements in the backend definition blocks are described in the following sections.

Some directives (`stickycookie` etc.) only have effect when Crossroads treats the network traffic as a stream of HTTP messages; i.e., when the service is declared with `type http`. In case of `type any`, the HTTP-specific directives have no effect.

4.4.1 `server` - Specifying the back end address

Description: Each back end must be identified by the network name (server name) where it is located. For example: `server 10.1.1.23`, or `server web.mydomain.org`. A TCP port specifier can follow the server name, as in `server web.mydomain.org:80`. **Note that** resolved host names can be cached by Crossroads. (The DNS cache timeout can be controlled using the invocation flag `-d`.)

Syntax: • `server servername`, where *servername* is a network name or IP address. In this case a separate `port` statement must be used to define the TCP port;

- `server servername:port`

Default: There is no default. This is a required setting.

4.4.2 port - Specifying a back end port

Description: Back ends must be known by their host name and a port. Both can be simultaneously specified in a `server` statement. When the `server` statement specifies a host name only, then a `port` statement must be used to specify the port.

Syntax: `port number`

Default: There is no default for the port. It must be specified either using `server` or using `port`.

4.4.3 verbosity - Controlling verbosity at the back end level

Description: Similar to `service` specifications, a `backend` can have its own verbosity (`on` or `off`). When `on`, traffic to and from this back end is reported.

Syntax: • `verbosity setting`, or
 • `verbose setting`, where `setting` is `true`, `yes` or `on`, or `false`, `no`, `off` to turn it off.

Default: `off`

4.4.4 retries - Specifying allowed failures

Description: Back ends that are 'flaky' or on a less reliable network can be marked as unavailable after not just one failure, but after e.g. three. You can use this configuration if you suspect that spurious errors cause otherwise 'good' back ends to be marked as unavailable, while they in fact still could be used.

Syntax: `retries number`; where `number` is the threshold of bad connections. Once exceeded, Crossroads will mark a back end as unavailable.

Default: 1; a back end is assumed to be unavailable after the first bad connection.

4.4.5 maxconnections - Limiting the connections to a back end

Description: The directive `maxconnections` limits the number of allowed connections to this client. Note that this directive can also occur on the level of a service block, in which case it limits the overall number of connections.

Futhermore note that this directive cannot be used when an external dispatcher is in effect. In such cases, the external dispatcher has full control over backend selection.

Syntax: `maxconnections number ;`

Default: 0; meaning no limit

4.4.6 weight - When a back end is more equal than others

Description: To influence how backends are selected, a backend can specify its 'weight' in the process. The higher the weight, the less likely a back end will be chosen. The default is 1. The weighing mechanism only applies to the dispatch modes `random`, `byconnections`, `bysize` and `byduration`. The weight is in fact a penalty factor. E.g., if backend A has weight 2 and backend B has weight 1, then backend B will be selected all the time, until its usage parameter is twice as large as the parameter of A. Think of it as a 'sluggishness' statement.

Syntax: *weight number*; the higher the number, the more 'sluggish' a back end is

Default: 1; all back ends have equal weight.

4.4.7 decay - Levelling out activity of a back end

Description: To make sure that a 'spike' of activity doesn't influence the perceived load of a back end forever, you may specify a certain decay. E.g, the statement `decay 10` makes sure that the load that crossroads computes for this back end (be it in seconds or in bytes) is decreased by 10% each time that **an other** back end is hit. Decays are not applied to the count of concurrent connections.

This means that when a given back end is hit, then its usage data of the transferred bytes and the connection duration are updated using the actual number of bytes and actual duration. However, when a different back end is hit, then the usage data are decreased by the specified decay.

Syntax: *decay number*, where *number* is a percentage that decreases the back end usage data when other back ends are hit.

Default: 0, meaning that no decay is applied to usage statistics.

4.4.8 state - Setting an initial back end state

Description: Using the `state` directive a back end can be 'primed' with an initial state. The keyword `state` can be followed by `available`, `unavailable` or `down`. Back ends marked `unavailable` are excluded as candidates, but are checked when a `revivinginterval` or a `checkinterval` is used. Back ends marked `down` are excluded and never re-checked.

Syntax: • *state specifier* ;
 • where *specifier* is one of `available`, `unavailable` or `down`

Default: `available`, meaning that a back end is a candidate for initial dispatching.

4.4.9 onstart, onend, onfail - Action Hooks

Description: The three directives `onstart`, `onend` and `onfail` can be specified to start system commands (external programs) when a connection to a back end starts, fails or ends:

- `onstart` commands will be run when Crossroads successfully connects to a back end, and starts servicing;
- `onend` commands will be run when a (previously established) connection stops;
- `onfail` commands will be run when Crossroads tries to contact a back end to serve a client, but the back end can't be reached.

The format is always *on`type` command*. The *command* is an external program, optionally followed by arguments. The command is expanded according to the following table:

- `%a` is the availability of the current back end, when a current back end is established. `%1a` is the availability of the first back end (0 when unavailable, 1 if available); `%2a` is the availability of the second back end, and so on.
- `%b` is the name of the current back end, when one is established. `%1b` is the name of the first back end, `%2b` of the second back end, and so on.
- `%e` is the count of seconds since start of epoch (January 1st 1970 GMT). `%60e` is the count since start of epoch plus 60, so this is a 1 minute offset into the future.
- `%g` is a "GMT-string" representation of the current time, in the format *monthname, day year hh:mm:ss*. This format is used in e.g. cookie expiry. `%600g` is the same representation but of a moment 600 seconds in the future (10 minutes).

- %h is the host name of the current back end. %1h is the host name of the first back end, %2h of the second back end, and so on.
- %l is the duration of the last successful connection in seconds, concerning the current back end. %1l is the duration of the connection to back end one, and so on. Note that %l refers only to the last successful connections. Unsuccessful connection attempts to back ends do not change the value.
- %p is the TCP port of the current back end. %1p is the TCP port of the first back end, %2p of the second back end, and so on.
- %r is the IP address of the connecting client.
- %s is the name of the current service that the client connected to.
- %t is the current local time in ANSI format, in YYYY-MM-DD/hh:mm:ss. %1800s is an ANSI stamp of a moment 1800 seconds in the future (half an hour).
- %T is the current GMT time in ANSI format. %10T offsets this 10 seconds into the future.
- %v is the Crossroads version.
- %w is the weight factor of the current back end. %1w is the weight factor of the first back end, etc..
- Any other character following a % sign is taken literally; e.g. %z is just a z.

Syntax: The syntax of the commands is as follows.

- `onstart commandline`
- `onend commandline`
- `onfail commandline`
- `onsuccess commandline`

Default: There is no default. Normally no external programs are run upon connection, success or failure of a back end.

4.4.10 `trafficlog` and `throughputlog` - Debugging and Performance Aids

Description: Two directives are available to log network traffic to files. They are `trafficlog` and `throughputlog`.

The `trafficlog` statement causes all traffic to be logged in hexadecimal format. Each line is prefixed by B or C, depending on whether the information was received from the back end or from the client.

The `throughputlog` statement writes shorthand transmissions to its log, accompanied by timings and the number of bytes each transmission was able to read or write in one chunk.

Syntax: The syntax is:

- `trafficlog filename ;`
- `throughputlog filename ;`

Default: none

4.4.11 `httptiminglog` - Timing debugging in HTTP mode

Description: The directive `httptiminglog` turns on logging of HTTP mode timing. There must be one argument, the filename where the timings are written to. Turning on this option will slow down processing, but may be helpful in finding out where delays are caused.

Syntax: The syntax is:

- `httptiminglog filename ;`
- Where *filename* is the output file, e.g. `/tmp/timings.log`

Default: none

4.4.12 stickycookie - Back end selection with an HTTP cookie

Description: The directive `stickycookie value` causes Crossroads to unpack clients' requests, to check for *value* in the cookies. When found, the message is routed to the back end having the appropriate `stickycookie` directive.

E.g., consider the following configuration:

```
service ... {  
    ...  
    backend one {  
        ...  
        stickycookie "BalancerID=first";  
    }  
    backend two {  
        ...  
        stickycookie "BalancerID=second";  
    }  
}
```

When clients' messages contain cookies named `BalancerID` with the value `first`, then such messages are routed to backend one. When the value is `second` then they are routed to the backend two.

There are basically two ways to provide such cookies to a browser. First, a back end can insert such a cookie into the HTTP response. E.g., the webserver of back end one might insert a cookie named `BalancerID`, having value `first`. Second, Crossroads can insert such cookies using a carefully crafted directive `addclientheader`.

Syntax: `stickycookie cookievalue`

Default: There is no default.

4.4.13 HTTP Header Modification Directives

Description: Crossroads understands the following header modification directives: `addclientheader`, `appendclientheader`, `setclientheader`, `addserverheader`, `appendserverheader`, `setserverheader`.

The directive names always consist of *ActionDestinationheader*, where:

- The action is `add`, `append` or `insert`.
 - Action `add` adds a header, even when headers with the same name already are present in an HTTP message. Adding headers is useful for e.g. `Set-Cookie` headers; a message may contain several of such headers.
 - Action `append` adds a header if it isn't present yet in an HTTP message. If such a header is already present, then the value is appended to the pre-existing header. This is useful for e.g. `Via` headers. Imagine an HTTP message with a header `Via: someproxy`. Then the directive `appendclientheader "Via: crossroads"` will rewrite the header to `Via: someproxy; crossroads`.
 - Action `set` overwrites headers with the same name; or adds a new header if no pre-existing is found. This is useful for e.g. `Host` headers.

- The destination is one of `client` or `server`. When the destination is `server`, then Crossroads will apply such directives to HTTP messages that originate from the browser and are being forwarded to back ends. When the destination is `client`, then Crossroads will apply such directives to backend responses that are shuttled to the browser.

The format of the directives is e.g. `addclientheader "X-Processed-By: Crossroads"`. The directives expect one argument; a string, consisting of a header name, a colon, and a header value. As usual, the directive must end with a semicolon.

The header value may contain one of the following formatting directives:

- `%a` is the availability of the current back end, when a current back end is established. `%1a` is the availability of the first back end (0 when unavailable, 1 if available); `%2a` is the availability of the second back end, and so on.
- `%b` is the name of the current back end, when one is established. `%1b` is the name of the first back end, `%2b` of the second back end, and so on.
- `%e` is the count of seconds since start of epoch (January 1st 1970 GMT). `%60e` is the count since start of epoch plus 60, so this is a 1 minute offset into the future.
- `%g` is a "GMT-string" representation of the current time, in the format *monthname, day year hh:mm:ss*. This format is used in e.g. cookie expiry. `%600g` is the same representation but of a moment 600 seconds in the future (10 minutes).
- `%h` is the host name of the current back end. `%1h` is the host name of the first back end, `%2h` of the second back end, and so on.
- `%l` is the duration of the last successful connection in seconds, concerning the current back end. `%1l` is the duration of the connection to back end one, and so on. Note that `%l` refers only to the last successful connections. Unsuccessful connection attempts to back ends do not change the value.
- `%p` is the TCP port of the current back end. `%1p` is the TCP port of the first back end, `%2p` of the second back end, and so on.
- `%r` is the IP address of the connecting client.
- `%s` is the name of the current service that the client connected to.
- `%t` is the current local time in ANSI format, in *YYYY-MM-DD/hh:mm:ss*. `%1800s` is an ANSI stamp of a moment 1800 seconds in the future (half an hour).
- `%T` is the current GMT time in ANSI format. `%10T` offsets this 10 seconds into the future.
- `%v` is the Crossroads version.
- `%w` is the weight factor of the current back end. `%1w` is the weight factor of the first back end, etc..
- Any other character following a `%` sign is taken literally; e.g. `%z` is just a `z`.

The following examples show common uses of header modifications.

Enforcing session stickiness: By combining `stickycookie` and `addclientheader`, HTTP session stickiness is enforced. Consider the following configuration:

```
service ... {
    ...
    backend one {
        ...
        addclientheader "Set-Cookie: BalancerID=first; path=/";
        stickycookie "BalancerID=first";
    }
    backend two {
        ...
        addclientheader "Set-Cookie: BalancerID=second; path=/";
```

```
        stickycookie "BalancerID=second";
    }
}
```

The first request of an HTTP session is balanced to either backend one or two. The server response is enriched using `addclientheader` with an appropriate cookie. A subsequent request from the same browser now has that cookie in place; and is therefore sent to the same back end where the its predecessors went.

The header which is sent to the client to inject a cookie, can furthermore be expanded to specify a timeout:

```
addclientheader "Set-Cookie: BalancerID=second; path=/; expires=%1800g";
```

The format specifier `%1800g` outputs a GMT-string date 1800 seconds in the future (half an hour from now).

Hiding the server software version: Many servers (e.g. Apache) advertize their version, as in `Server: Apache 1.27`. This potentially provides information to attackers. The following configuration hides such information:

```
service ... {
    ...
    backend one {
        ...
        setclientheader "Server: WWW-Server";
    }
}
```

Informing the server of the clients' IP address: Since Crossroads sits 'in the middle' between a client and a back end, the back end perceives Crossroads as its client. The following sends the true clients' IP address to the server, in a header `X-Real-IP`:

```
service ... {
    ...
    backend one {
        ...
        setserverheader "X-Real-IP: %r";
    }
}
```

Keep-Alive Downgrading: The directives `setclientheader` and `setserverheader` also play a key role in downgrading Keep-Alive connections to 'single-shot'. E.g., the following configuration makes sure that no Keep-Alive connections occur.

```
service ... {
    ...
    backend one {
        ...
        setserverheader "Connection: close";
        setclientheader "Connection: close";
    }
}
```

- Syntax:**
- `addclientheader Headername: headervalue` to add a header in the traffic towards the client, even when another header *Headername* exists;
 - `appendclientheader Headername: headervalue` to append *headervalue* to an existing header *Headername* in the traffic towards the client, or to add the whole header altogether;
 - `setclientheader Headername: headervalue` to overwrite an existing header in the traffic towards the client, or to add such a header;
 - `addserverheader Headername: headervalue` to add a header in the traffic towards the server, even when another header *Headername* exists;

- `appendserverheader` *Headername: headervalue* to append *headervalue* to an existing header *Headername* in the traffic towards the server, or to add the whole header altogether;
- `setserverheader` *Headername: headervalue* to overwrite an existing header in the traffic towards the server, or to add such a header.

Default: There is no default.

5 Tips, Tricks and Random Remarks

The following sections elaborate on the directives as described in section 4 to illustrate how crossroads works and to help you achieve the "optimal" balancing configuration.

5.1 Configuration examples

5.1.1 A load balancer for three webserver back ends

The following configuration example binds crossroads to port 80 of the current server, and distributes the load over three back ends. This configuration shows most of the possible settings.

```
service www {
    /* We don't need session stickyness. */
    type any;

    /* Port on which we'll listen in this service: required. */
    port 8000;

    /* What IP address should this service listen? Default is 'any'.
     * Alternatively you can state an explicit IP address, such as
     * 127.0.0.1; that would bind the service only to 'localhost'. */
    bindto any;

    /* Verbose reporting or not. Default is off. */
    verbosity on;

    /* Dispatching mode, or: How to select a back end for an incoming
     * request. Possible values:
     * roundrobin: just the next back end in line
     * random: like roundrobin, but at random to make things more
     *          confusing. Probably only good for testing.
     * bysize: The backend that transferred the least nr of bytes
     *          is the next in line. As a modifier you can say e.g.
     *          bysize over 10, meaning that the 10 last connections will
     *          be used to compute the transfer size, instead of all
     *          transfers.
     * byduration: The backend that was active for the shortest time
     *              is the next in line. As a modifier you can say e.g.
     *              byduration of 10 to compute over the last 10 connections.
     * byconnections: The back end with the least active connections
     *                 is the next in line.
     * byorder: The first available back end is always taken.
     */
    dispatchmode byduration over 5;
```

```
/* Interval at which we'll check whether a temporarily unavailable
 * backend has woken up.
 */
revivinginterval 5;

/* TCP backlog of connections. Default is 0 (no backlog, one
 * connection may be active).
 */
backlog 5;

/* For status reporting: a shared memory key. Default is the same
 * as the port number, OR-ed by a magic number.
 */
shmkey 8000;

/* This controls when crossroads should consider a connection as
 * finished even when the TCP sockets weren't closed. This is to
 * avoid hanging connections that don't do anything. NOTE THAT when
 * crossroads cuts off a connection due to timeout exceed, this is
 * not marked as a failure, but as a success. Default is 0: no timeout.
 */
connectiontimeout 300;

/* The max number of allowed client connections. When present, connections
 * won't be accepted if the max is about to be exceeded. When
 * absent, all connections will be accepted, which might be misused
 * for a DOS attack.
 */
maxconnections 300;

/* Now let's define a couple of back ends. Number 1: */
backend www_backend_1 {
    /* The server and its port, the minimum configuration. */
    server httpserver1;
    port 9010;
    /* The 'decay' of usage data of this back end. Only relevant
     * when the whole service has 'dispatchmode bysize' or
     * 'byduration'. The number is a percentage by which the usage
     * parameter is decreased upon each connection of an other back
     * end.
     */
    decay 10;

    /* To see what's happening in /var/log/messages: */
    verbosity on;
}

/* The second one: */
backend www_backend_2 {
    /* Server and port */
    server httpserver2;
    port 9011;

    /* Verbosity of reporting when this back end is active */
```

```
    verbosity on;

    /* Decay */
    decay 10;

    /* This back end is twice as weak as the first one */
    weight 2;

    /* Event triggers for system commands upon succesful activation
    * and upon failure.
    */
    onsuccess echo 'success on backend 2' | mail root;
    onfailure echo 'failure on backend 2' | mail root;
}

/* And yet another one.. this time we will dump the traffic
* to a trace file. Furthermore we don't want more than 10 concurrent
* connections here. Note that there's also a total maxconnections for the
* whole service.
*/
backend www_backend_3 {
    server httpserver3;
    verbosity on;
    port 9000;
    verbosity on;
    decay 10;
    trafficlog /tmp/backend.3.log;
    maxconnections 10;
}
}
```

5.1.2 An HTTP forwarder when travelling

As another example, here's my `/etc/crossroads.conf` that I use on my Unix laptop. The problem that I face is that I need many HTTP proxy configurations (at home, at customers' sites and so on) but I'm too lazy to reconfigure browsers all the time.

Here's how it used to be before crossroads:

- At home, I would surf through a squid proxy on my local machine. The browser proxy setting is then `http://localhost:3128`.
- Sometimes I start up an SSH tunnel to our offices. The tunnel has a local port 3129, and connects to a squid proxy on our e-tunity server. Hence, the browser proxy is then `http://localhost:3129`.
- At a customer's location I need the proxy `http://10.120.34.113:8080`, because they have configured it so.
- And in yet other instances, I use a HTTP diagnostic tool Charles⁴ that sits between browser and website and shows me what's happening. I run `charles` on my own machine and it listens to port 8888, behaving like a proxy. The browser configuration for the proxy is then `http://localhost:8888`.

Here's how it works with a crossroads configuration:

⁴<http://www.xk72.com/charles>

- I have configured my browsers to use `http://localhost:8080` as the proxy. For all situations.
- I use the following crossroads configuration, and let crossroads figure out which proxy backend works, and which doesn't. Note two particularities:
 - The statement `dispatchmode byorder`. This makes sure that once crossroads determines which backend works, it will stick to it. This usage of crossroads doesn't need to balance over more than one back end.
 - The statement `bindto 127.0.0.1` makes sure that requests from other interfaces than loopback won't get serviced.

```
service HttpProxy {
    port 8080;
    bindto 127.0.0.1;
    verbosity on;
    dispatchmode byorder;
    revivinginterval 15;

    backend Charles {
        server localhost:8888;
        verbosity on;
    }

    backend CustomerProxy {
        server 10.120.34.113:8080;
        verbosity on;
    }

    backend SshTunnel {
        server localhost:3129;
    }

    backend LocalSquid {
        server localhost:3128;
    }
}
```

As a final note, the commandline argument `tell` can be used to influence crossroad's own detection mechanism of back end availability detection. E.g., if in the above example the back ends `SshTunnel` and `LocalSquid` are both active, then crossroads `tell httpproxy sshtunnel down` will 'take down' the back end `SshTunnel` – and will automatically cause crossroads to switch to `LocalSquid`.

5.1.3 SSH login with enforced idle logout

The following example shows how crossroads 'throttles' SSH logins. Connections are accepted on port 22 (the normal SSH port) and forwarded to the actual SSH daemon which is running on port 2222.

Note the usage of the `connectiontimeout` directive. This makes sure that users are logged out after 10 minutes of inactivity. Note also the `maxconnections` setting, this makes sure that no more than 10 concurrent logins occur.

```
service Ssh {
    port 22;
    backlog 5;
```

```
maxconnections 10;
connectiontimeout 600;
backend TrueSshDaemon {
    server localhost:2222;
}
}
```

5.2 How back ends are selected in load balancing

In order to tune your load balancing, you'll need to understand how crossroads computes usage, how weighing works, and so on. In this section we'll focus on the dispatching modes `bysize`, `byduration` and `byconnections` only. The other dispatching types are self-explanatory.

5.2.1 Bysize, byduration or byconnections?

As stated before, crossroads doesn't know 'what a service does' and how to judge whether a given back end is very busy or not. You must therefore give the right hints:

- In general, a service which is CPU bound, will be more busy when it takes longer to process a request. The dispatch mode `byduration` is appropriate here.
- In contrast, a service which is filesystem bound, will be more busy when more data are transferred. The dispatch mode `bysize` is appropriate.
- The dispatch mode `byduration` can also be used when network latency is an issue. E.g., if your balancer has back ends that are geographically distributed, then `byduration` would be a good way to select best available back ends.
- Furthermore it is noteworthy that `dispatchmode byduration` is not usable for interactive processes such as SSH logins. Idle time of a login adds to the duration, while causing (almost) no load. Mode `byduration` should only be used for automated processes that don't wait for user interaction (e.g., SOAP calls and other HTTP requests).
- As a last remark, the dispatching mode `byconnections` can be used if you don't have other clues for load estimations.

E.g., consider a database connection. What's heavier on the back end, time-consuming connections, or connections where loads of bytes are transferred? Well, that depends. A tough `select` query that joins multiple tables can be very heavy on the back end, though the response set can be quite small - and hence the number of transferred bytes. That would suggest dispatching by duration. However, `byduration` balancing doesn't represent the true world, when interactive connections can occur where users have an idle TCP connection to the database: this consumes time, but no bytes (see the SSH login example above). In this case, the dispatch mode `byconnections` may be your best bet.

5.2.2 Averaging size and duration

The configuration statement `dispatchmode bysize` or `byduration` allows an optional modifier `over number`, where the stated number represents a connection count. When this modifier is present, then crossroads will use a moving average over the last *n* connections to compute duration and size figures.

In the real world you'll always want this modifier. E.g., consider two back ends that are running for years now, and one of them is suddenly overloaded and very busy (it experiences a 'spike' in activity). When the `over` modifier is absent, then the sudden load will hardly show up in the usage figures - it will flatten out due to the large usage figures already stored in the years of service.

In contrast, when e.g. `over 3` is in effect, then a sudden load does show up - because it highly contributes to the average of three connections.

5.2.3 Specifying decays

Decays are also only relevant when crossroads computes the 'next best back end' by size (bytes) or duration (seconds). E.g., imagine two back ends A and B, both averaged over say 3 connections.

Now when back end A is suddenly hit by a spike, its average would go up accordingly. But the back end would never again be used, unless B also received a similar spike, because A's 'usage data' over its last three connections would forever be larger than B's data.

For that reason, you should in real situations probably always specify a decay, so that the backend selection algorithm recovers from spikes. Note that the usage data of the back end where a decay is specified, decay when **other** back ends are hit. The decay parameter is like specifying how fast your body regenerates when someone else does the work.

The below configuration illustrates this:

```
/* Definition of the service */
service soap {
    /* Local TCP port */
    port 8080;

    /* We'll select back ends by the processing
     * duration
     */
    dispatchmode byduration over 3;

    /* First back end: */
    backend A {
        /* Back end IP address and port */
        server 10.1.1.1:8080;

        /* When this back end is NOT hit because
         * the other one was less busy, then the
         * usage parameters decay 10% per connection
         */
        decay 10;
    }

    /* Second back end: */
    backend B {
        server 10.1.1.2:8080;
        decay 10;
    }
}
```

5.2.4 Adjusting the weights

The back end modifier `weight` is useful in situations where your back ends differ in respect to performance. E.g., your back ends may be geographically distributed, and you know that a given back end is difficult to reach and often experiences network lag.

Or you may have one primary back end, a system with a fast CPU and enough memory, and a small fall-back back end, with a slow CPU and short on memory. In that case you know in advance that the second back end should be used only rarely. Most requests should go to the big server, up to a certain load.

In such cases you will know in advance that the best performing back ends should be selected the most often. Here's where the `weight` statement comes in: you can simply increase the weight of the back ends with the least performance, so that they are selected less frequently.

E.g., consider the following configuration:

```
service soap {
  port 8080;
  dispatchmode byduration over 3;
  backend A {
    server 10.1.1.1:8080;
    decay 20;
  }
  backend B {
    server 10.1.1.2:8080;
    weight 2;
    decay 10;
  }
  backend C {
    server 10.1.1.3:8080;
    weight 4;
    decay 5;
  }
}
```

This will cause crossroads to select back ends by the processing time, averaging over the last three connections. However, backend B will kick in only when its usage is half of the usage of A (back end B is probably only half as fast as A). Backend C will kick in only when its usage is a quarter of the usage of A, which is half of the usage of B (back end C is probably very weak, and just a fall-back system incase both A and B crash). Note also that A's usage data decay much faster than B's and C's: we're assuming that this big server recovers quicker than its smaller siblings.

5.3 Periodic probes and wake up calls

Crossroads has two methods of periodic back end verifications:

- The first method only checks unavailable back ends. It is configured using the directive `revivinginterval`. The way `revivinginterval` works is as follows:
 1. Crossroads determines that a back end is unavailable. This happens when a new network connection to the back end fails, or when an existing connection suddenly dies.
 2. The back end is the periodically checked to see if it has woken up yet.
- The second method periodically checks all back ends, available ones and unavailable ones. This is configured using `checkinterval`. During each check, back ends can be marked as available or as unavailable.

5.3.1 Syntax

The syntax of both verifications is:

```
method interval seconds [ externalhandler program arguments ]
```

In this syntax, the method is `reviving` or `check`. The seconds specifier defines the delay between consecutive checks.

When no specification `externalhandler program arguments` is given, then Crossroads runs the verification by trying to connect to the back end. A simple successful connection suffices to cause Crossroads to consider the back end available. However, when the `externalhandler` specifier is given, then Crossroads runs the specified external program. This program must exit

with status 0 to inform Crossroads that the back end is available. All other exit statuses will mark the back end as unavailable.

The *arguments* are expanded according to the following table:

- %a is the availability of the current back end, when a current back end is established. %1a is the availability of the first back end (0 when unavailable, 1 if available); %2a is the availability of the second back end, and so on.
- %b is the name of the current back end, when one is established. %1b is the name of the first back end, %2b of the second back end, and so on.
- %e is the count of seconds since start of epoch (January 1st 1970 GMT). %60e is the count since start of epoch plus 60, so this is a 1 minute offset into the future.
- %g is a "GMT-string" representation of the current time, in the format *monthname, day year hh:mm:ss*. This format is used in e.g. cookie expiry. %600g is the same representation but of a moment 600 seconds in the future (10 minutes).
- %h is the host name of the current back end. %1h is the host name of the first back end, %2h of the second back end, and so on.
- %l is the duration of the last successful connection in seconds, concerning the current back end. %1l is the duration of the connection to back end one, and so on. Note that %l refers only to the last successful connections. Unsuccessful connection attempts to back ends do not change the value.
- %p is the TCP port of the current back end. %1p is the TCP port of the first back end, %2p of the second back end, and so on.
- %r is the IP address of the connecting client.
- %s is the name of the current service that the client connected to.
- %t is the current local time in ANSI format, in *YYYY-MM-DD/hh:mm:ss*. %1800s is an ANSI stamp of a moment 1800 seconds in the future (half an hour).
- %T is the current GMT time in ANSI format. %10T offsets this 10 seconds into the future.
- %v is the Crossroads version.
- %w is the weight factor of the current back end. %1w is the weight factor of the first back end, etc..
- Any other character following a % sign is taken literally; e.g. %z is just a z.

5.3.2 Security Considerations

When `externalhandler` is in effect, Crossroads spawns an external process. For security reasons, you may want to run this process under a restricted user account.

The directive `useraccount` can be used to accomplish this.

5.3.3 An example

The following configuration balances SMTP requests to two back ends. The connectivity is checked by retrieving output from each back end. The back end is available when the standard greeting of a mail server is seen; this greeting must contain the word SMTP. During each check the probe is terminated by sending `quit`.

```
/* Use either DEBUG on or off */
#define DEBUG off

/* SMTP balancer */
service smtp {
    /* Standard stuff */
```



```
port 25;
verbosity DEBUG;
/* Check back ends every 30 seconds. User 'nobody'
 * will run the external handler. */
useraccount nobody;
checkinterval 30
    externalhandler "echo quit | netcat -wl %h %p | grep SMTP";

/* Two back ends to handle mail traffic. */
backend mailone {
    server smtp1.local.network:25;
    verbosity DEBUG;
}
backend second {
    server smtp2.local.network:25;
    verbosity DEBUG;
}
}
```

5.4 Throttling the number of concurrent connections

If you suspect that your service may occasionally receive ‘spikes’ of activity⁵, then it might be a good idea to protect your service by specifying a maximum number of concurrent connections. This protection can be specified on two levels:

On the service level a statement like `maxconnections 100;` states that the service as a whole will never service more than 100 concurrent connections. This means that all your back ends and the crossroads balancer itself will be protected from being overloaded.

On the back end level a statement like `maxconnections 10;` states that this particular back end will never have more than 10 concurrent connections; regardless of the overall setting on the service level. This means that this particular back end will be protected from being overloaded (regardless of what other back ends may experience).

The `maxconnections` statement, combined with a back end selection algorithm, allows very fine granularity. The `maxconnections` statement on the back end level is like a hand brake: even when you specify a back end algorithm that would protect a given back end from being used too much, a situation may occur where that back end is about to be hit. A `maxconnections` statement on the level of that back may then protect it.

5.5 TCP Session Stickiness

If you need to make sure that a client that once gets dispatched to a given back end keeps re-visiting the back end, then Crossroads offers the dispatch mode `byclientip`. This mode will only work when each client is seen by Crossroads with its own specific IP address; ie., this method won’t work when clients reach Crossroads through a masquerading firewall (in which case all clients would be seen as having the firewall’s IP address).

The `dispatchmode byclientip` works as follows:

- The client’s IP address is taken in its string representation and ‘hashed’ into a number.
- The number is brought back to the number of available back ends (using a modulo-operation).
- The result defines the back end of choice.

⁵which you should always assume

If the preferred back end is unavailable, then the action that Crossroads takes is to dispatch as if `byconnections`: of the available back ends, the one with the least connections is taken.

5.6 HTTP Session Stickiness

This section focuses on HTTP session stickiness. This term refers to the ability of a balancer to route a conversation between browser and a backend farm with web servers always to the same back end. In other words: once a back end is selected by the balancer, it will remain the back end of choice, even for subsequent connections.

5.6.1 Don't use stickiness!

The rule of thumb as far as the balancer is concerned, is: **Do not use HTTP session stickiness unless you really have to**. Enabling session stickiness hampers failover, balancing and performance:

- Failover is hampered because during the session, the balancer has to assign new connections to the same back end that was selected at the start of a session. If the back end suddenly goes 'down', then the session will most likely crash. (Actually, when a back end becomes unreachable in the middle of a session, Crossroads will assign a new back end to that session. This will most likely result in a malfunction of the underlying application.)
- Balancing is hampered because at the start of the session, the balancer has selected the next-best back end. But during the session, that back end may well become overloaded. The balancer however must continue to send the requests there.
- Performance is hampered because crossroads needs to 'unpack' messages as they are passed to and fro. That's because crossroads needs to check the HTTP headers in the messages for persistence cookies.

There is a number of measures that you can take to avoid using session stickiness. E.g., session data can be 'shared' between web back ends. PHP offers functionality to store session data in a database, so that all PHP applications have access to these data. Application servers such as Websphere can be configured to replicate session data between nodes.

5.6.2 But if you must..

If you really need stickiness, think first whether you might use TCP stickiness (using the client's IP address to dispatch). If you can, then this is the preferred method, since Crossroads won't have to unpack TCP streams. Below is a short configuration example:

```
service www {
    port 80;
    type any;
    revivinginterval 15;
    dispatchmode byclientip;

    backend one {
        server 10.1.1.100:80;
    }

    backend two {
        server 10.1.1.101:80;
    }
}
```

However, if you **must** use HTTP-base session stickiness, then proceed as follows:

- At the level of a `service` description, set the type to `http`.
- Furthermore, set the `headerinspection` to `shallow` (unless of course you also want to modify other HTTP headers, see section 5.8).
- At the level of each back end description, configure the `stickycookie` and a `addclientheader` directives.

Once crossroads sees that, it will examine each HTTP message that it shuttles between client and back end:

- If there is no persistence cookie in the HTTP headers of a client's request, then the message must be the first one and a new session should be established. Crossroads selects an appropriate back end, sends the message to that back end, catches the reply, and inserts a `Set-Cookie` directive.
- If there is a persistence cookie in the HTTP headers of a client's request, then the request is part of an already established session. Crossroads analyzes the cookie and forwards the request to the appropriate back end.

Below is a short example of a configuration.

```
service www {
    port 80;
    type http;
    headerinspection shallow;
    revivinginterval 15;
    dispatchmode byconnections;

    backend one {
        server 10.1.1.100:80;
        stickycookie XRID=100;
        addclientheader "Set-Cookie: XRID=100; Path=/";
        /* or: XRID=100; Path=/; Expires=%600g
        * This would make the session cookie expire in 600 sec (10 mins)
        */
    }

    backend two {
        server 10.1.1.101:80;
        stickycookie XRID=101;
        addclientheader "Set-Cookie: XRID=101; Path=/";
    }
}
```

Note how the cookie names and values in the directives `stickycookie` and `addclientheader` match. That is obviously a prerequisite for stickiness.

5.7 Passing the client's IP address

Since Crossroads just shuttles bytes to and fro, meta-information of network connections is lost. As far as the back ends are concerned, their connections originate at the Crossroads junction. For example, standard Apache access logs will show the IP address of Crossroads.

In order to compensate for this, Crossroads can insert a special header in HTTP connections, to inform the back end of the original client's IP address. In order to enable this, the Crossroads configuration must state the following:

- The service type must be `http`, and not `any`;

- Make sure that the `headerinspection` is deep (or that there is no `headerinspection` statement, since `deep` is the default, see section 5.8);
- In the back end definition, the following statement must occur:
`addserverheader "X-Real-IP: %r";`
You are of course free to choose the header name; the here used `X-Real-IP` is a common name for this purpose.

After this, HTTP traffic that arrives at the back ends has a new header: `X-Real-IP`, holding the client's IP address. **Note that** once the type is set to `http`, Crossroads' performance will be hampered – all passing messages will have to be unpacked and analyzed.

5.7.1 Sample Crossroads configuration

The below sample configuration shows two HTTP back ends that receive the client's IP address:

```
service www {
    port 80;
    type http;
    revivinginterval 5;
    dispatchmode roundrobin;

    backend one {
        server 10.1.1.100:80;
        addserverheader "X-Real-IP: %r";
    }

    backend two {
        server 10.1.1.200:80;
        addserverheader "X-Real-IP: %r";
    }
}
```

5.7.2 Sample Apache configuration

The method by which each back end analyzes the header `X-Real-IP` will obviously be different per server implementations. However, a common method with the Apache webserver is to log the client's IP address into the access log.

Often this is accomplished using the log format `custom`, defined as follows:

```
LogFormat "%h %l %u %t %D \"%r\" %>s %b" common
CustomLog logs/access_log common
```

The first line defines the format `common`, with the remote host specified by `%h`. The second line sends access information to a log file `logs/access_log`, using the previously defined format `common`.

Furtunately, Apache's `LogFormat` allows one to log contents of headers. By replacing the `%h` with `%{X-Real-IP}i`, the desired information is sent to the log. Therefore, normally you can simply redefine the `common` format to

```
LogFormat "%{X-Real-IP}i %l %u %t %D \"%r\" %>s %b" common
```

5.8 Deep or shallow HTTP header inspection

The service-level directive `headerinspection` defines which HTTP headers Crossroads will analyze. Often, several HTTP requests and responses will be served over one network link. Browsers and servers will try to keep a TCP link open so that it may be re-used; this is a

measure to increase efficiency and to shorten load times of pages. E.g., a typical HTML page will require a style sheet and a few images - and these can be retrieved over the same link that originally served the HTML page.

Re-using the TCP link occurs more often than not. It is the default in HTTP/1.1 (unless `Connection: close` is specified as one of the HTTP headers). The older HTTP protocol, HTTP/1.0, by default passes just one request and response over a TCP link, after which the link is closed (unless `Connection: keep-alive` is specified as one of the HTTP headers).

If you define your service as `type http`, then by default Crossroads will inspect all HTTP headers that it sees: not only the first browser request and server answer, but also subsequent requests and answers that travel over the same link. This is called 'deep inspection mode', in which Crossroads applies directives such as `addclientheader` to all header blocks. Inspecting the full TCP stream to catch header blocks can be resource consuming. You can optionally make Crossroads's resource consumption 'lighter' by instructing it to inspect only the first HTTP header block, and to simply pass-through over all subsequent information (which might well include next header blocks of a re-used TCP connection) This is done using the directive `headerinspection`:

```
service web {
    type http;
    headerinspection shallow;
    backend a { .... }
    backend b { .... }
}
```

The situations where `shallow` mode can be used, depends on what you need to do:

- `shallow` mode can be used when the inspection and modification of the first HTTP header block suffices. For example, HTTP session stickiness is a good example:

```
service web {
    type http;
    headerinspection shallow;
    backend a {
        server 10.1.1.1:80;
        stickycookie "BalancerID=1";
        addclientheader "Set-Cookie: BalancerID=1";
    }
    .
    . other back ends are defined here
    .
}
```

In this case, Crossroads will inspect only the first header block that the client sends for the presence of a cookie `BalancerID`. If the cookie has value 1, then the request will be routed to back end a. Similarly, Crossroads will inspect only the first header block that the server sends, and will inject a `Set-Cookie` instruction.

In this example, deep mode is not necessary because Crossroads will use the first header that the client sends to route the information to a given back end. If more than one HTTP transactions follow over the same TCP link, then by definition the link will go to the same back end - even without inspecting all HTTP headers that follow the first ones.

- `deep` mode is necessary in situations where all header blocks must be inspected and modified. E.g., if you want to hide your HTTP server identification, you might use:

```
service web {
    type http;
    headerinspection deep;
```

```
    backend a {
        server 10.1.1.1:80;
        setclientheader "Server: MyWebServer";
    }
    .
    . other back ends are defined here
    .
}
```

Here, all HTTP header blocks that come from the server will be parsed, and `Server` headers will be overwritten. In a similar vein, if you want to pass the client's IP address, you will also need `deep` mode.

In these examples, `shallow` mode is not usable, because the outbound header modifications should apply to all headers of a given series. Imagine that one would use `shallow` mode here: then, in a series of 5 HTTP transactions that pass over the same TCP link, only the first transaction would hide the HTTP server signature. All subsequent transactions would still show the HTTP server signature to the world.

5.9 Debugging network traffic

Incase the traffic between client and backend must be debugged, the statement `trafficlog filename` can be issued. This causes the traffic to be dumped in hexadecimal format to the stated filename.

Traffic sent by the client is prefixed by a **C**, traffic sent by the back end is prefixed by a **B**. Below is a sample traffic dump of a browser trying to get a HTML page. The server replies that the page was not modified.

```
C 0000  47 45 54 20 68 74 74 70 3a 2f 2f 77 77 77 2e 63 GET http://www.c
C 0010  73 2e 68 65 6c 73 69 6e 6b 69 2e 66 69 2f 6c 69 s.helsinki.fi/li
C 0020  6e 75 78 2f 6c 69 6e 75 78 2d 6b 65 72 6e 65 6c nux/linux-kernel
C 0030  2f 32 30 30 31 2d 34 37 2f 30 34 31 37 2e 68 74 /2001-47/0417.ht
C 0040  6d 6c 20 48 54 54 50 2f 31 2e 31 0d 0a 43 6f 6e ml HTTP/1.1..Con
C 0050  6e 65 63 74 69 6f 6e 3a 20 63 6c 6f 73 65 0d 0a nection: close..
.
. etcetera
.
B 0000  48 54 54 50 2f 31 2e 30 20 33 30 34 20 4e 6f 74 HTTP/1.0 304 Not
B 0010  20 4d 6f 64 69 66 69 65 64 0d 0a 44 61 74 65 3a Modified..Date:
B 0020  20 54 75 65 2c 20 31 32 20 4a 75 6c 20 32 30 30 Tue, 12 Jul 200
B 0030  35 20 30 39 3a 34 39 3a 34 37 20 47 4d 54 0d 0a 5 09:49:47 GMT..
B 0040  43 6f 6e 74 65 6e 74 2d 54 79 70 65 3a 20 74 65 Content-Type: te
B 0050  78 74 2f 68 74 6d 6c 3b 20 63 68 61 72 73 65 74 xt/html; charset
.
. etcetera
.
```

Turning on traffic dumps will *significantly* slow down crossroads.

Besides `trafficlog`, there is also a directive `throughputlog`. This directive also takes one argument, a filename. The file is appended, and the following information is logged:

- The process ID of the crossroads image that serves the TCP connection;
- The time of the request, in seconds and microseconds since start of the run;
- A **C** when the request originated at the client, or **B** when the request originated at the back end;
- The first 100 bytes of the request.

As an example, consider the following (the lines are shortened for brevity and prefixed by line numbers for clarity):

```
1 0027566 0.000000 C 462 GET http://public.e-tunity.com/ ...
2 0027566 0.052974 B 394 HTTP/1.1 200 OK..Via: 1.1 tinyp ...
3 0027566 0.053413 B 1073 <html>.. <head>.. <titl ...
4 0027566 0.053589 B 1448 1> </td>.. </tr>.. </table ...
5 0027566 0.065679 B 725 for more. info.... <li> To ...
```

This tells us that:

- Line 1: PID 27566 served a request that originated at the client (C). The corresponding timing is 0 seconds, which means that this is the start of the run. The request was 462 bytes long and started with GET http://public...
- Line 2: 0.052974 seconds later, the backend (B) replied with HTTP/1.1 200 OK.
- Lines 3, 4 and 5: The backend (B) proceeded by sending chunks.

The buffer sizes (462, 394 and so on) also tell us that the network sends relatively small chunks. In this case we might save some processing memory by setting the TCP buffer size to say 2K, instead of the default 5K. In this situation, most of the default 5K buffers will be unused. Setting the TCP buffer size to 2K can be done using the flag `-B 2048` or by using the statement `tcpbuffersize 2048`; in an options block of the configuration.

It is also worth while remembering that the start time of a C request is the time that crossroads sees the activity. Any latency between the true client and crossroads is obviously not included. This is illustrated by the below simple ASCII art:

```
client ---->---->---->---->*crossroads =====>=====>
                                                    \
                                                    back end
                                                    /
client ----<----<----<----<----< crossroads =====<=====<=====<
```

This simple picture shows a typical request that originates at a client, travels to crossroads, and is relayed via the back end. The C entry in a throughput log is the time when crossroads sees the request, indicated by an asterisk. The B entries are the times that it takes the back end to answer, indicated by == style lines. Therefore, the true roundtrip time will be longer than the number of seconds that are logged in the throughput log: the latency between client and crossroads isn't included in that measurement.

Summarizing, the throughput times and buffer sizes of a client-back end connection can be analyzed using the directive `throughputlog`. In a real-world analysis, you'd probably want to write up a script to analyze the output and to compute round trip times. Such scripts are not (yet) included in Crossroads.

5.10 IP filtering: Limiting Access by Client IP Address

5.10.1 General Examples

The directives `allowfrom`, `denyfrom`, `allowfile` and `denyfile` can be used to instruct Crossroads to specifically allow access by using a "whitelist" of IP addresses, or to specifically deny access by using a "blacklist". E.g., the following configuration allows access to service `webproxy` only to `localhost`:

```
service webproxy {
    port 8000;
    allowfrom 127.0.0.1;
```

```
backend one {  
    .  
    . Back end definitions occur here  
    .  
}  
.  
. Other back ends or other service directives  
. may occur here  
.
```

In this example there is a "whitelist" having only one entry: IP address 127.0.0.1, or *localhost*. (Incidentally, the same behaviour could be accomplished by stating *bindto 127.0.0.1*, in which case Crossroads would only listen to the local network device.)

In the same vein, the directive `allowfrom 127.0.0.1 192.168.1/24` would allow access to *localhost* and to all IP addresses that start with 192.168.1. The specifier `192.168.1/24` states that there are three network bytes (192, 168 and 1), and 24 bits (or 3 bytes) are relevant; so that the fourth network byte doesn't matter.

5.10.2 Using External Files

The directives `allowfile` and `denyfile` allow you to specify IP addresses in external files. The Crossroads configuration states e.g. `allowfile /tmp/allow.txt`, and the IP addresses are then in `/tmp/allow.txt`. The format of `/tmp/allow.txt` is as follows:

- The specifications follow again *p.q.r.s/mask*, where p, q, r and s are network bytes which can be left out on the right hand side when the mask allows it;
- The specifications must be separated by white space (spaces, tabs or newlines).

E.g., the following is a valid example of an external specification file:

```
127.0.0.1  
192.168.1/24  
10/8
```

When external files are in effect, then the signal `SIGHUP` (1) causes Crossroads to reload the external file. E.g., while Crossroads is running, you may edit `/tmp/allow.txt`, and then issue `killall -1 crossroads`. The new contents of `/tmp/allow.txt` will be reloaded.

5.10.3 Mixing Directives

Crossroads allows to mix all directives in one service description. However, some mixes are less meaningful than others. It's up to you to take this into account.

The following rules apply:

- Blacklisting and whitelisting can be used together. When combined, the blacklist will always be interpreted first. E.g., consider the following directives:

```
allowfrom 192.168.1/24  
denyfrom 192.168.1.100
```

Given the fact that the deny list is checked first, client 192.168.1.100 won't be able to access Crossroads. Then the allow list will be checked, stating that all clients whose IP address starts with 192.168.1 may connect. The effect will be that e.g., client 192.168.1.1 may connect, 192.168.1.2 may connect too, 192.168.1.100 will be blocked, and 10.1.1.1 will be blocked as well.

Now consider the following directives:


```
allowfrom 192.168.1.100 127.0.0.1
denyfrom 192.168.1/24
```

This will first of all deny access to all IP addresses that start with 192.168.1. So the rule that allows 192.168.1.100 won't ever be effective. The net result will be that access will be granted to 127.0.0.1.

- Blacklisting or whitelisting can be left out. A list is considered empty when no appropriate directives occur in `/etc/crossroads.conf`, or when the directive points to an empty or non-existent external file.
- Using `*from` and `*file` statements is allowed, but doesn't make sense. E.g., the following configuration sample is such a case:

```
allowfrom 127.0.0.1 192.168.1/24
allowfile /tmp/allow.txt
```

There is a technical reason for this. Once Crossroads processes the `allowfile` directive, then the whole whitelist is cleared (thereby removing the entries 127.0.0.1 and 192.168.1/24), and new entries are reloaded from the file. The net result is that the `allowfrom` specification is overruled.

Crossroads only performs syntactic checking of the configuration. Some of the above samples are syntactically correct, but make no semantic sense: Crossroads doesn't warn for such situations.

5.11 Using an external program to dispatch

As mentioned before, Crossroads supports several built-in dispatch modes. However, you are always free to hook-in your own dispatch mode that determines the next back end using your own specific algorithm. This section explains how to do it.

5.11.1 Configuring the external handler

First, the `dispatchmode` statement needs to inform Crossroads that an external program will do the job. The syntax is: `dispatchmode externalhandler program arguments`. The *program* must point to an executable program that will be started by Crossroads. The specifier *arguments* can be anything you want; those will be the arguments to your dispatch helper. You use the following special format specifiers in the argument list:

- `%a` is the availability of the current back end, when a current back end is established. `%1a` is the availability of the first back end (0 when unavailable, 1 if available); `%2a` is the availability of the second back end, and so on.
- `%b` is the name of the current back end, when one is established. `%1b` is the name of the first back end, `%2b` of the second back end, and so on.
- `%e` is the count of seconds since start of epoch (January 1st 1970 GMT). `%60e` is the count since start of epoch plus 60, so this is a 1 minute offset into the future.
- `%g` is a "GMT-string" representation of the current time, in the format *monthname, day year hh:mm:ss*. This format is used in e.g. cookie expiry. `%600g` is the same representation but of a moment 600 seconds in the future (10 minutes).
- `%h` is the host name of the current back end. `%1h` is the host name of the first back end, `%2h` of the second back end, and so on.
- `%l` is the duration of the last successful connection in seconds, concerning the current back end. `%1l` is the duration of the connection to back end one, and so on. Note that `%l` refers only to the last successful connections. Unsuccessful connection attempts to back ends do not change the value.

- `%p` is the TCP port of the current back end. `%1p` is the TCP port of the first back end, `%2p` of the second back end, and so on.
- `%r` is the IP address of the connecting client.
- `%s` is the name of the current service that the client connected to.
- `%t` is the current local time in ANSI format, in `YYYY-MM-DD/hh:mm:ss`. `%1800s` is an ANSI stamp of a moment 1800 seconds in the future (half an hour).
- `%T` is the current GMT time in ANSI format. `%10T` offsets this 10 seconds into the future.
- `%v` is the Crossroads version.
- `%w` is the weight factor of the current back end. `%1w` is the weight factor of the first back end, etc..
- Any other character following a `%` sign is taken literally; e.g. `%z` is just a `z`.

Note that the format specifiers such as `%b` don't make sense in the phase in which an external handler is called, since there is no current back end yet (the job of the handler is to supply one, so at the time of calling, `%b` is undefined).

5.11.2 Writing the external handler

The external handler is activated using the arguments that are specified in `/etc/crossroads.conf`. The external handler can do whatever it wants, but ultimately, it must write a back end name on its *stdout*. Crossroads reads this, and if the back end is available, uses that back end for the connection.

5.11.3 Examples of external handlers

This section shows some examples of Crossroads configurations vs. external handlers. The sample handlers that are shown here, are also included in the Crossroads distribution, under the directory `etc/`. Also note that the examples shown here are just quick-and-dirty Perl scripts, meant to illustrate only. Your applications may need other external handlers, but you can use the shown scripts as a starting point.

Round-robin dispatching This example is trivial in the sense that round-robin dispatching is already built into Crossroads, so that using an external handler for this purpose only slows down Crossroads. However, it's a good starting example.

The Crossroads configuration is shown below:

```
service test {
    port 8001;
    verbosity on;
    revivinginterval 5;

    dispatchmode externalhandler
        /usr/local/src/crossroads/etc/dispatcher-roundrobin
        %1b %1a %2b %2a;

    backend testone {
        server localhost:3128;
        verbosity on;
    }
    backend testtwo {
        server localhost:3128;
        verbosity on;
    }
}
```

The relevant `dispatchmode` statement invokes the external program `dispatcher-roundrobin` with four arguments: the name of the first back end (`testone`), its availability (0 or 1), the name of the second back end (`testtwo`) and its availability (0 or 1).

The external handler, which is also included in the Crossroads distribution, is shown below. It is a Perl script.

```
#!/usr/bin/perl

use strict;

# Example of a round-robin external dispatcher. This is totally
# superfluous, Crossroads has this on-board; if you use the external
# program for determining round-robin dispatching, then you'll only
# slow things down. This script is just meant as an example.

# Globals / configuration
# -----
my $log = '/tmp/exthandler.log';    # Debug log, set to /dev/null to suppress
my $statefile = '/tmp/rr.last';    # Where we keep the last used

# Logging
# -----
sub msg {
    return if ($log eq '/dev/null' or $log eq '');
    open (my $of, ">>$log") or return;
    print $of (scalar(localtime()), ' ', @_);
}

# Read the last used back end
# -----
sub readlast() {
    my $ret;

    if (open (my $if, $statefile)) {
        $ret = <$if>;
        chomp ($ret);
        close ($if);
        msg ("Last used back end: $ret\n");
        return ($ret);
    }
    msg ("No last-used back end (yet)\n");
    return (undef);
}

# Write back the last used back end, reply to Crossroads and stop
# -----
sub reply ($) {
    my $last = shift;

    if (open (my $of, ">$statefile")) {
        print $of ("$last\n");
    }
    print ("$last\n");
    exit (0);
}
```

```
# Main starts here
# -----

# Collect the cmdline arguments. We expect pairs of backend-name /
# backend-availability, and we'll store only the available ones.
msg ("Dispatch request received\n");
my @backend;
for (my $i = 0; $i <= $#ARGV; $i += 2) {
    push (@backend, $ARGV[$i]) if ($ARGV[$i + 1]);
}
msg ("Available back ends: @backend\n");

# Let's see what the last one is. If none found, then we return the
# first available back end. Otherwise we need to go thru the list of
# back ends, and return the next one in line.
my $last = readlast();
if ($last eq '') {
    msg ("Returning first available back end $backend[0]\n");
    reply ($backend[0]);
}

# There was a last back end. Try to match it in the list,
# then return the next-in-line.
for (my $i = 0; $i < $#backend; $i++) {
    if ($last eq $backend[$i]) {
        msg ("Returning next back end ", $backend[$i + 1], "\n");
        reply ($backend[$i + 1]);
    }
}

# No luck.. run back to the first one.
msg ("Returning first back end $backend[0]\n");
reply ($backend[0]);
```

The working of the script is basically as follows:

- The argument list is scanned. Back ends that are available are collected in an array `@backend`.
- The script queries a state file `/tmp/rr.last`. If a back end name occurs there, then the next back end is looked up in `@backend` and returned to Crossroads. If no last back is unknown or can't be matched, then the first available back end (first element of `@backend`) is returned to Crossroads.
- Informing Crossroads is done via the subroutine `reply()`. This code writes the selected back end to file `/tmp/rr.last` (for future usage) and prints the back end name to *stdout*.
- The script logs its actions to a file `/tmp/exthandler.log`. This log file can be inspected for the script's actions.

Dispatching by the client IP address The following example shows a useful real-life situation that illustrates how dispatching by client IP address works. **Note that** as of Crossroads 1.31, `dispatchmode byclientip` is implemented – so that the below description is somewhat superfluous. The code snippets however can help you in modelling your own specific dispatch modes, aided by external helpers. (Incidentally, the `dispatchmode byclientip`

was modeled after the shown script. The functionality proved so useful that it was embedded into Crossroads.)

Our hypothetical dispatching situation is as follows:

- Crossroads is used as a single-address point to forward Remote Desktop requests to a farm of Windows systems, where users can work via remote access;
- However, users may stop their session, and when they re-connect, they expect to be sent to the Windows system that they had worked on previously;
- Client PC's have their distinct IP addresses, which distinguish them.
- Of four windows systems, two are large servers, and two are small ones. We'll want to assign large servers to clients when we have a choice.

The requirements resemble session stickiness in HTTP, except that the remote desktop protocol doesn't support stickiness. This situation is a perfect example of how an external handler can help: the potential delay due to the calling of an external handler won't even be noticed. Remote Desktop is a network service where the connection time isn't critical; users won't notice a slightly larger connection time due to the fact that Crossroads invokes an external program. We expect only a few (albeit lengthy) TCP connections.

The approach to the solution of this problem uses several external program hooks:

- An external dispatcher handler will be responsible for suggesting a back end, given a client IP and given the current timestamp. This handler will consult an internal administration to see whether the stated IP address should re-use a back end, or to determine which back end is free for usage.
- An external hook `onstart` will be responsible for updating the internal administration; i.e., to flag a back end as 'occupied'.
- The external hooks `onfailure` and `onend` will be responsible for flagging a back end as 'free' again; i.e., for erasing any previous information that states that the back end was occupied.

The Crossroads configuration is shown below. Only four Windows back ends are shown. Each back end is configured on a given IP address, port 3389, and is limited to one concurrent connection (otherwise a new user might 'steal' a running desktop session).

```
service rdp {
    port 3389;
    revivinginterval 5;

    /* rdp-helper dispatch IP STAMP ... will suggest a back end to use,
     * arguments are for all back ends: name, availability, weight */
    dispatchmode externalhandler
        /usr/local/src/crossroads/etc/rdp-helper dispatch %r %e
        %1b %1a %1w
        %2b %2a %2w
        %3b %3a %3w
        %4b %4a %4w;

    backend win1 {
        server 10.1.1.1:3389;
        maxconnections 1;
        /* rdp-helper start IP STAMP BACKEND will log the actual start
         * of a connection;
         * rdp-helper end IP will log the ending of a connection */
        onstart /usr/local/src/crossroads/etc/rdp-helper start %r %e %b;
        onend /usr/local/src/crossroads/etc/rdp-helper end %r;
```

```
        onfail /usr/local/src/crossroads/etc/rdp-helper end %r;
    }
    backend win2 {
        server 10.1.1.2:3389;
        maxconnections 1;
        onstart /usr/local/src/crossroads/etc/rdp-helper start %r %e %b;
        onend /usr/local/src/crossroads/etc/rdp-helper end %r;
        onfail /usr/local/src/crossroads/etc/rdp-helper end %r;
    }
    backend win3 {
        server 10.1.1.3:3389;
        maxconnections 1;
        weight 2;
        onstart /usr/local/src/crossroads/etc/rdp-helper start %r %e %b;
        onend /usr/local/src/crossroads/etc/rdp-helper end %r;
        onfail /usr/local/src/crossroads/etc/rdp-helper end %r;
    }
    backend win4 {
        server 10.1.1.4:3389;
        maxconnections 1;
        weight 3;
        onstart /usr/local/src/crossroads/etc/rdp-helper start %r %e %b;
        onend /usr/local/src/crossroads/etc/rdp-helper end %r;
        onfail /usr/local/src/crossroads/etc/rdp-helper end %r;
    }
}
```

Depending on the dispatcher stage, the external handler `rdp-helper` is invoked in different ways:

During dispatching the helper is called to suggest a back end. The arguments are an action indicator `dispatch`, the client's IP address, the timestamp, and four triplets that represent back ends: per back end its name, its availability, and its weight. The purpose of the helper is to tell Crossroads which back end to use.

During connection start the helper will be invoked to inform it of the start of a connection, given a client IP address.

When a connection terminates the helper will be invoked to inform it that the connection has ended.

Here's the external handler as Perl script. It uses the module `GDBM_File` which most likely will not be part of standard Perl distributions, but can be added using CPAN. (Alternatively, any other database module can be used.)

```
#!/usr/bin/perl

use strict;
use GDBM_File;

# Global variables and configuration
# -----
my $log = '/tmp/exthandler.log'; # Debug log, set to /dev/null to suppress
my $cdb = '/tmp/client.db';     # GDBM database of clients
my %db;                         # .. and memory representation of it
my $timeout = 24*60*60;         # Timeout of a connection in secs
```

```
# Logging
# -----
sub msg {
    return if ($log eq '/dev/null' or $log eq '');
    open (my $of, ">>$log") or return;
    print $of (scalar(localtime()), ' ', @_);
    close ($of);
}

# Reply a back end to the caller and stop processing.
# -----
sub reply ($) {
    my $b = shift;
    msg ("Suggesting $b to Crossroads.\n");
    print ("$b\n");
    exit (0);
}

# Is a value in an array
# -----
sub inarray {
    my $val = shift;
    for my $other (@_) {
        return (1) if ($other eq $val);
    }
    return (0);
}

# A connection is starting
# -----
sub start {
    my ($ip, $stamp, $backend) = @_;
    msg ("Logging START of connection for IP $ip on stamp $stamp, ",
        "back end $backend\n");
    $db{$ip} = "$backend:$stamp";
}

# A connection has ended
# -----
sub end {
    my $ip = shift;
    msg ("Logging END of connection for IP $ip\n");
    $db{$ip} = undef;
}

# Request to determine a back end
# -----
sub dispatch {
    my $ip = shift;
    my $stamp = shift;

    msg ("Request to dispatch IP $ip on stamp $stamp\n");

    # Read the next arguments. They are triplets of
```

```
# backend-name / availability / weight. Store if the back end is
# available.
my (@backends, @weights);
for (my $i = 0; $i < $#_; $i += 3) {
    if ($_[ $i + 1] != 0) {
        push (@backends, $_[ $i]);
        push (@weights, $_[ $i + 2]);
        msg ("Candidate back end: $_[ $i] with weight ", $_[ $i + 2], "\n");
    }
}

# See if this is a reconnect by a previously seen client IP. We'll
# treat this as a reconnect if the timeout wasn't yet exceeded.
if ($db{$ip} ne '') {
    my ($last_backend, $last_stamp) = split (/:/, $db{$ip});
    msg ("IP $ip had last connected on $last_stamp to $last_backend\n");
    if ($stamp < $last_stamp + $timeout) {
        msg ("Timeout not yet exceeded, this may be a reconnect\n");
        # We'll allow a reconnect only if the stated last_backend is
        # free (sanity check).
        if (inarray ($last_backend, @backends)) {
            msg ("Last back end $last_backend is available, ",
                "letting through\n");
            reply ($last_backend);
        } else {
            msg ("Last used back end isn't free, suggesting a new one\n");
        }
    } else {
        msg ("Timeout exceeded, suggesting a new back end\n");
    }
} else {
    msg ("No previous connection data, suggesting a new back end\n");
}

my $bestweight = -1;
my $bestbackend;
for (my $i = 0; $i <= $#weights; $i++) {
    if ($bestweight == -1 or $bestweight > $weights[$i]) {
        $bestweight = $weights[$i];
        $bestbackend = $backends[$i];
    }
}

msg ("Best back end: $bestbackend (given weight $bestweight)\n");
reply ($bestbackend);
}

# Main starts here
# -----
msg ("Start of run, attaching GDBM database '$cdb'\n");
tie (%db, 'GDBM_File', $cdb, &GDBM_WRCREAT, 0600);

# The first argument must be an action 'dispatch', 'start' or 'end'.
# Depending on the action, we do stuff.
```



```
my $action = shift (@ARGV);
if ($action eq 'dispatch') {
    dispatch (@ARGV);
} elsif ($action eq 'start') {
    start (@ARGV);
} elsif ($action eq 'end') {
    end (@ARGV);
} else {
    print STDERR ("Usage: rdp-helper {dispatch|start|end} args\n");
    exit (1);
}
```

5.12 Linux and ip_conntrack_max

The kernel value of `ip_conntrack_max` is important for routers and balancers under Linux. Basically it's the maximum number of tracked connections. Felix A.W.O. describes the following situation:

- Crossroads seems to mark back ends as unavailable, while in fact nothing is afoot.
- This happens under heavy load.
- In `/var/log/messages` one may see the message: `kernel: ip_conntrack: table full, dropping packet.`

The reason for Crossroads's behavior is that the kernel refuses to build up a requested network connection. For Crossroads, this looks just as a non-responding back end. Crossroads therefore marks the back end as unavailable.

The solution is as follows:

- Try `cat /proc/sys/net/ipv4/ip_conntrack_max` to see the current value.
- Add something to the shown value (e.g., multiply by two), and inform the kernel of the new value, using `echo new-value > /proc/sys/net/ipv4/ip_conntrack_max`
- Make sure that the same step occurs somewhere in your boot sequence as well, or that the new value is stated in a configuration file such as `/etc/sysctl.conf`.

The value for *new-value* is something that you'll have to figure out yourself. Note however that each count will cause the kernel to reserve 350 bytes. So if you set `ip_conntrack_max` to 100.000, then you're already taking 33.3Mb off the total available memory.

5.13 Marking back ends as bad after more than one try

Crossroads allows you to specify on a per-back end basis how many retries are needed before a back end is considered unavailable. The default is just one, meaning that after one failed connection, Crossroads will mark a back end as unavailable (the back end may be 'revived', if you use `revivinginterval` or `checkinterval`).

Increasing the number is specified using the keyword `retries`. The following configuration defines two back ends; the one on the IP address 5.6.7.8 is somehow 'flaky', and Crossroads should try connecting 3 times before crossing it off:

```
service www {
    port 80;
    backend plugh {
        server 1.2.3.4:80;
    }
    backend xyzy {
        server 5.6.7.8:80;
    }
}
```

```
        retries 3;
    }
}
```

There may be several reasons for increasing the retries number:

- The network connections to the server may spuriously hamper, but such rare errors don't mean that the back end server is down.
- The back end server is a 'slow starter' and can't handle spikes very well. E.g., it may be a webserver which starts with only one daemon; extra capacity is added as network connections arrive, but adding capacity take a little time.

Whatever the reason, the keyword `retries` might be of help here. This keyword should however be used carefully: Crossroads will retry connecting with a small one-second delay in between. A high `retries` number means also lots of one-second delays, in which time a client is kept waiting.

5.14 Using the Web Interface `crossroads-mgr`

The mini-webserver `crossroads-mgr` provides an intuitive web interface to the state of Crossroads. Once started, an administrator may view the state of the balancer, and may influence the state of all back ends.

As an example, a procedure is described here where a given back end is gracefully taken out of service. Below is a sample Crossroads configuration, which distributes requests over three back ends:

```
service www {
    port 80;
    type http;
    backend one {
        server 10.1.1.1:80;
        stickycookie "BalancerID=one";
        addclientheader "Set-Cookie: BalancerID=one";
    }
    backend two {
        server 10.1.1.1:80;
        stickycookie "BalancerID=two";
        addclientheader "Set-Cookie: BalancerID=two";
    }
    backend three {
        server 10.1.1.1:80;
        stickycookie "BalancerID=three";
        addclientheader "Set-Cookie: BalancerID=three";
    }
}
```

In order to take a given back end gracefully offline (say back end `two`), either of the following procedures can be followed:

1. Back end `two` is taken down via the commandline, using `crossroads tell www two down`. Next we wait for say one hour to let existing sessions die out. After that, the back end can be turned off.
2. Back end `two` is taken down via the web interface, by switching the status from `available` to `down`. Next we wait for say one our, and turn off the back end.

The web interface doesn't offer extra functionality over the command line tools; but all information is available at one glance, and accessible without a shell access to the balancer.

5.14.1 Starting crossroads-mgr

The basic command to start `crossroads-mgr` is

```
crossroads-mgr start portnumber
```

where the port number specifies to which TCP port the manager will listen. There is however one important security aspect that needs attention. Unless specified otherwise, anyone who points their browser to the balancer and the given port, can view back end states and even change them. This may be a too lax policy.

The web interface daemon has two methods to limit access: a listening address can be specified via the command line, and basic authentication (username / password protection) can be turned on.

Specifying a listening address is done using the flag `-a`. E.g., after

```
crossroads-mgr -a 127.0.0.1 start 10000
```

the manager is started to listen to port 10000, but only on the address 127.0.0.1 which is localhost. Requests from other addresses will not be served.

Enforcing basic authentication is turned on using two command line flags:

- Flag `-b` turns on basic authentication to view the status.
- Flag `-B` turns on basic authentication to change back end states.

Both flags must be followed by the required credentials. The most simple way to state the credentials, is to postfix the flag with the required user name, a colon, and the required password. E.g., `-Buser:secret` says that anyone who tries to change back end states, must supply the user name `user` and the password `secret`.

In this example the full invocation would be e.g.:

```
crossroads-mgr -Buser:secret start 1000
```

The disadvantage is here that the credentials are visible for anyone who has shell access to the balancer. A process overview, generated with say `ps ax |grep crossroads-mgr`, would show the required username and password. In order to avoid such a leak, `crossroads-mgr` can be started as follows:

```
crossroads-mgr -BPROMPT start 1000
```

The 'magic' word `PROMPT` instructs `crossroads-mgr` to read the username and password from *stdin*. The invocation can be further scripted, using something like:

```
echo user:secret |crossroads-mgr -BPROMPT start 1000
```

The same trick can be used with the flag `-b`. When both flags are present, and both 'magic' words `PROMPT` occur, then `crossroads-mgr` will first ask for the credentials of the 'viewer', and next for the credentials of the 'modifier' (even when flag `-B` is stated first). So the following example starts `crossroads-mgr` and requires the user name `viewer`, with password `showme` to view the status, and it requires the user name `modifier`, with password `changeit` to change states:

```
echo -e 'viewer:showme\nmodifier:changeit' |  
crossroads-mgr -bPROMPT -BPROMPT start 1000
```

5.15 Rendering Crossroads' status in a web page

The Crossroads flag `-x` causes the status output to be generated as an XML document. This format can be nicely used to render the output as an HTML page.

E.g., the following sample shows how Crossroads reports its status in XML format (the actual XML structure that Crossroads outputs may be different from this example, depending on the version):

```
<?xml version="1.0" encoding="UTF-8"?>
<status>
  <service id="1" name="smtp">
    <connections>1</connections>
    <lastbackend>0</lastbackend>
    <backend id="0" name="first">
      <availability id="0">available</availability>
      <clients>0</clients>
      <failures>0</failures>
      <connections>2</connections>
      <duration sec="0.559882">0.56s</duration>
      <throughput bytes="3564">3.48Kb</throughput>
    </backend>
    <backend id="1" name="second">
      <availability id="0">available</availability>
      <clients>0</clients>
      <failures>0</failures>
      <connections>2</connections>
      <duration sec="23.7636">23.76s</duration>
      <throughput bytes="9055">8.84Kb</throughput>
    </backend>
  </service>
</status>
```

A custom-made XSLT transformation stylesheet can be used to convert this to any output format - and also to HTML. Such a style sheet is included in the Crossroads distribution as `etc/xml-status-to-html.xslt`. The sheet is lengthy, and is therefore not included in this document. (You're welcome to modify it to suit your specific needs. If you have cool tips, send them along to me and I'll include them in the next distribution!)

If you want to show this output in a webpage which is generated on demand by a web-server, then you might run into the following problem. The status reporter (`crossroads -x status`) must be able to access the shared memory segment of the running Crossroads instance. By default, the shared memory is protected for the user that started Crossroads, which will often not be the user who runs the webserver. Under the auspices of the webserver user, `crossroads status` might abort with a message: "ERROR: Cannot get shared memory for service *name*, key *number*: Permission denied."

The solution for this problem is to make the shared memory access somewhat more liberal. There are basically two options:

- Start Crossroads with the flag `-m 0666`, which makes the shared memory segment available to all. The octal number 0666 works just like a file permission setting under Unix. Now, any user can run `crossroads status`.
- If you want to make the access to Crossroads' shared memory somewhat stricter, then make the user who starts Crossroads and the user who runs the webserver member of a new Unix group. Start Crossroads with the flag `-m 0664`. Now, users belonging to the same group as the Crossroads starter can run `crossroads status`.

5.16 Crossroads and DNS caching

The option `-d` allows you to control Crossroads' built in DNS caching mechanism. Most often you will not need to use this: DNS lookups in Crossroads occur only to find back ends; and back ends are usually 'near' to the balancer.

You might want or need to use DNS caching if:

- Back end servers are specified as hostnames, and not as IP addresses;
- DNS resolving of those host names is perceptibly slow.

You can test DNS resolving from the command line using e.g. the commands `nslookup` and `host`.

If DNS resolving is an issue, then you can specify the flag `-d nsec` on the command line while invoking Crossroads. This instructs Crossroads to use its DNS cache to store results. Each result is stored for up to `nsec` seconds - after that, a new request for the back end will lead to a new DNS lookup.

5.17 Managing a Pool of Virtual Back Ends

Crossroads can be used as a central dispatcher for computing on demand. E.g., imagine a situation where a web service exists on one or two servers, with the option of adding new servers in case of a high load. The "other" servers, besides the first two ones, are a virtual pool - the pool size and the IP addresses may not be known in advance.

The following Crossroads' features are used in this example:

- The back ends of Crossroads can have an initial state. For the two real servers which are already present this would be `up` (the default). For the non-present ones this would be `down`: Crossroads excludes back ends that are down from the set of candidate workers, and does not check their presence using `revivalinterval` or `checkinterval`.
- The back ends must be configured to have a server address and port, but this can be later reconfigured runtime. E.g., the virtual pool servers would be initially configured on `localhost:80`, and later the addresses would be filled in.

A sample Crossroads configuration is shown below. It defines 20 back ends: two are always present, and 18 are in the virtual pool.

```
service vpool {
    port 80;
    revivinginterval 3;
    dispatchmode byconnections;

    /* The two back ends that are always present: */
    backend real_01 { server 10.1.1.1:80; }
    backend real_02 { server 10.1.1.2:80; }

    /* The virtual pool: */
    backend virt_01 { server localhost:80; state down; }
    backend virt_02 { server localhost:80; state down; }
    backend virt_03 { server localhost:80; state down; }
    backend virt_04 { server localhost:80; state down; }
    backend virt_05 { server localhost:80; state down; }
    backend virt_06 { server localhost:80; state down; }
    backend virt_07 { server localhost:80; state down; }
    backend virt_08 { server localhost:80; state down; }
    backend virt_09 { server localhost:80; state down; }
    backend virt_10 { server localhost:80; state down; }
```

```
backend virt_11 { server localhost:80; state down; }
backend virt_12 { server localhost:80; state down; }
backend virt_13 { server localhost:80; state down; }
backend virt_14 { server localhost:80; state down; }
backend virt_15 { server localhost:80; state down; }
backend virt_16 { server localhost:80; state down; }
backend virt_17 { server localhost:80; state down; }
backend virt_18 { server localhost:80; state down; }

}
```

New back ends can be enabled using `crossroads tell`. E.g., to enable back end `virt_01` on IP address `10.100.1.1:80`, the commands are:

```
crossroads tell vpool virt_01 server 10.100.1.1:80
crossroads tell vpool virt_01 up
```

To disable a back end, its state is reset to down. The IP address doesn't even have to be erased:

```
crossroads tell vpool virt_01 down
```

A script to monitor the total number of connections, and to add a back end or to remove one, is left to the reader. The output of `crossroads -x status` can be very helpful here: it reports on the states of all back ends, their connections, the total number of available or down back ends, etc..

6 Benchmarking

This section shows how `crossroads` affects the transmitting of HTML data when used as an intermediate 'station' through which all data travels.

6.1 Benchmark 1: Accessing a proxy via `crossroads` or directly

The benchmark was run on a system where the following was varied:

1. A website was recursively spidered through a local squid proxy. The spidering was repeated 10 times, the total was recorded.
2. `Crossroads` was placed in front of the squid proxy, and the website was again recursively spidered. Again, the spidering was repeated 10 times and the total was recorded.

The `crossroads` configuration of the second alternative is shown below:

```
service HttpProxy {
    port 8080;
    verbosity on;
    backend LocalSquid {
        server 127.0.0.1;
        port 3128;
        verbosity on;
    }
}
```

6.1.1 Results

The results of this test are that crossroads causes a negligible delay, if it is statistically relevant at all. Without crossroads, the timing results are:

```
real 0m8.146s
user 0m0.130s
sys 0m0.253s
```

When using crossroads as a middle station, the results are:

```
real 0m9.481s
user 0m0.141s
sys 0m0.230s
```

6.1.2 Discussion

The above shown results are quite favorable to crossroads. However, one should know that situations will exist where crossroads leans towards the 'worst case' scenario, causing up to 50% delay.

E.g., imagine a test where a `wget` command retrieves a HTML document from an Apache server on `localhost`. Now we have (almost) no overhead due to network throttling, host-name lookups and so on. When this test would be run either with or without crossroads in between, then theoretically, crossroads would cause a much larger delay, because it has to read from the server, and then write the same information to `wget`. Each read/write occurs twice when crossroads sits in between.

This worst case scenario will however (fortunately) occur only very seldom in the real world:

- Normally network issues, such as the above mentioned host name lookups or throughput restrictions, will add significantly to the duration of a request. The 'twice as many' read/writes caused by crossroads are then relatively irrelevant.
- Normally a significant amount of time will be spent in a back end, due to processing (e.g., when calling a servlet on a back end). Again, this processing time will weigh much heavier than the multiple read/writes.

6.2 Benchmark 2: Crossroads versus Linux Virtual Server (LVS)

LVS is a kernel-based balancer that acts like a masquerading firewall: TCP packets that arrive at the balancer are sent to one of the configured back ends. LVS has the advantage over crossroads that there is no stop-and-go in the transmission; in contrast, crossroads needs to send data via an internal buffer. Crossroads has the advantage that it offers instantaneous failover because it tries to contact the back end for upon each new TCP connection; in contrast, LVS isn't aware of downtime of back ends (unless one implements an external heartbeat). Also, crossroads offers more complex balancing than LVS.

6.2.1 Environment

On the balancer, LVS was run on port 80, its forwarding set up for two equally weighted back ends, using `ipvsadm`:

```
ipvsadm -a -t 192.168.1.250:http -r 10.1.1.100:http -m -w 1
ipvsadm -a -t 192.168.1.250:http -r 10.1.1.101:http -m -w 1
```

Crossroads was run on port 81. The configuration file is shown below:

```
service http {
    port 81;
    dispatchmode roundrobin;
    revivinginterval 5;
    backend one {
        server 10.1.1.100;
        port 80;
    }
    backend two {
        server 10.1.1.101;
        port 80;
    }
}
```

6.2.2 Tests and results

In the first test, ports 80 and 81 on the balancer were 'bombed' with 50 concurrent clients, each requesting a small page 50 times. The following timings where measured:

- How long it takes to establish a connection;
- How long it takes to retrieve the page.

The results of this test were:

- On average, each client took 0.12 seconds to connect to LVS, and each page was retrieved in 0.14 seconds;
- On average, each client took 0.11 seconds to connect to crossroads, and each page was retrieved in 0.13 seconds.

In this setup there seems to be no difference between the performance of LVS and crossroads!

In a second test, the size of the retrieved page was varied from 2.000 to 2.000.000 bytes. This test was taken to see whether crossroads would show performance degradation when transferring larger amounts of data.

For each page size, 30 concurrent clients were started, that retrieved the page 50 times. Again, the connect times and processing times where recorded.

The results of the total time (connect time + retrieval time) are shown in the below table:

Bytes	LVS timing	Crossroads timing
2000	0.130741688	0.12739582
20000	0.490916224	0.50376901
200000	3.799440328	4.33125273
2000000	45.25090855	45.9600728

Again, the results show that crossroads performs just as effectively as LVS, even with large data chunks!

7 Compiling and Installing

7.1 Prerequisites

The creation of crossroads requires:

- Standard Unix tools, such as `sed`, `awk`;
- Perl (5.00 or better). The web interface `crossroads-mgr` requires specific modules;
- A POSIX-compliant C compiler;

- Support for SYSV IPC, networking and so on.

Basically a Linux or Apple MacOSX box will do nicely. To compile and install crossroads, follow these steps.

7.2 Compiling and installing

- Obtain the source distribution. It can be found on <http://crossroads.e-tunity.com>. The distribution comes as an archive `crossroads-type.tar.gz`, where *type* is *stable* or *devel*.
- Unpack the archive in a sources directory using `tar xzf crossroads-X.YY.tar.gz`. The contents spill into a subdirectory `crossroads-X.YY/`.
- Change-dir into the directory.
- Next, edit `etc/Makefile.def` and verify that all compilation settings are to your likings. The settings are explained in the file. **Note that** the default distribution of `Makefile.def` is suited for Linux or Apple MacOSX systems. On other Unices, or on non-Unix systems, you must particularly pay attention to `SET_PROC_TITLE_BY...`. When in doubt, set the `SET_PROC_TITLE...` settings to 0. Crossroads will work nevertheless, but it won't show nice titles in `ps` listings. Also there's a macro `EXTRA_LIBS` to add linkage flags (an example for a Solaris build is included).
- Now crossroads is ready for compilation. Do a `make local` followed by `make install`. The latter step may have to be done by the user `root` if the `BINDIR` setting of `etc/Makefile.def` points to a root-owned directory.
- The manual pages on the basic usage of Crossroads and on the layout of the configuration file are installed as well, if a suitable `man/` directory is present under the installation prefix directory. If these pages are not installed, then you can always copy `doc/crossroads.1` to a suitable `man/man1/` directory, and `doc/crossroads.conf.7` to a suitable `man/man7/` directory. After this, `man crossroads` and `man crossroads.conf` would show the appropriate manual pages.
- The full documentation doesn't install in this process. If you want to install the documentation, then proceed as follows:
 - Optionally, `cp doc/crossroads.html htmldirectory/`; where *htmldirectory* is the destination directory for your HTML manuals;
 - Optionally, `cp doc/crossroads.pdf pdfdirectory/`; where *pdfdirectory* is the destination directory for your PDF manuals.
- In order to use the web interface `crossroads-mgr`, make sure that the following Perl modules are available: `HTTP::Daemon`, `Getopt::Std`, `MIME::Base64`. When in doubt, start the commandline tool `cpan` and enter `install HTTP::Daemon`, which will install the named module if necessary. You can repeat this for all listed modules.

7.3 Configuring crossroads

Now that the binary is available on your system, you need to create a suitable `/etc/crossroads.conf`. Use this manual to get started.

Once you have the configuration ready, start crossroads with `crossroads start`. Test the availability of your services and back ends. Monitor how crossroads is doing with:

- In one terminal, run the script:

```
while [ 1 ] ; do
    tput clear
    crossroads status
    sleep 3
done
```

Note that depending on your system you might need `sleep 3s`, i.e., with an `s` appended.

- In another terminal, run:

```
while [ 1 ] ; do
    tput clear
    ps ax | grep crossroads | grep -v grep
    sleep 3
done
```

Note that depending on your system you might need `ps -ef` instead of `ps ax`.

- In yet another terminal, run `tail -f /var/log/messages` (supply the appropriate system log file if `/var/log/messages` doesn't work for you).

Now thoroughly test the availability of your back ends through crossroads. The status display will show an updated view of which back ends are selected and how busy they are. The process list will show which crossroads daemons are running. Finally, the tailing of `/var/log/messages` shows what's going on – especially if you have `verbosity true` statements in the configuration.

7.4 A boot script

Finally, you may want to create a boot-time startup script. The exact procedure depends on the used Unix flavor.

7.4.1 SysV Style Startup

On SysV style systems, there's a startup script directory `/etc/init.d` where bootscripts for all utilities are located. You may have the `chkconfig` utility to automate the task of inserting scripts into the boot sequence, but otherwise the steps will resemble the following.

- Create a script `crossroads` in `/etc/init.d` similar to the following:

```
#!/bin/sh
/usr/local/bin/crossroads -v $1
/usr/local/bin/crossroads-mgr $1 1000
```

`init` will supply the right value for `$1`: "start" during startup, and "stop" during shutdown. Note that the stated directory `/usr/local/bin` must correspond with the installation path.

In the first line, The flag `-v` causes the startup to be more 'verbose'. However, once daemonized, the verbosity of Crossroads is controlled by the appropriate statements in the configuration.

The second line starts the web interface `crossroads-mgr` on TCP port 1000. Access to the web interface is free for all; add your own flags `-b` or `-B` to enforce basic authentication (see the manual page of `crossroads-mgr` and section 5.14).

- Determine your 'runlevel': usually 3 when your system is running in text-mode only, or 5 when you are using a graphical interface. If your runlevel is 3, then:

```
root> cd /etc/rc.d/rc3.d
root> ln -s /etc/init.d/crossroads S99crossroads
root> ln -s /etc/init.d/crossroads K99crossroads
```

This creates startup (`S*`) and stop (`K*`) links that will be run when the system enters or leaves a given runlevel.

If your runlevel is 5, then the right `cd` command is to `/etc/rc.d/rc5.d`. Alternatively, you can create the symlinks in both runlevel directories.

7.4.2 BSD Style Startup

On BSD style systems, daemons are booted directly from `/etc/rc` and related scripts. In case you have a file `/etc/rc.local`, edit it, and add the statement:

```
/usr/local/bin/crossroads start  
/usr/local/bin/crossroads-mgr start 1000
```

If your BSD system lacks `/etc/rc.local`, then you may need to start Crossroads from `/etc/rc`. Your mileage may vary.

7.4.3 Linux-related

When using Crossroads on Linux, the following may be relevant:

- Crossroads supports flag `-p`, which causes Crossroads not to modify the process name. In the absence of the flag, Crossroads alters its process name to something like `crossroads - Service myservice: listening`.

Flag `-p` may be handy if you want to suppress this behavior. The process name is then `crossroads-daemon`. Some boot scripts use the `/proc` filesystem to find PID's by filename; in those cases, flag `-p` can be used.